

Streamlet: Textbook Streamlined Blockchains

Benjamin Chan
Cornell University

Joint work with Elaine Shi

cbr.stanford.edu/sbc20/

“Simplifying Consensus”

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1. Modeling consensus (5min)
2. Motivating simplicity as a goal (a few seconds)
3. Our protocol (20min)

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Goal: walk away

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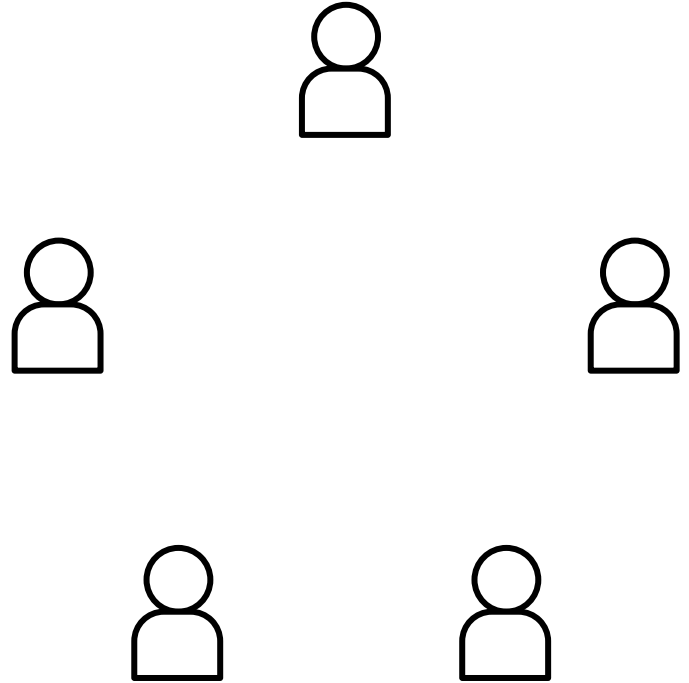
Goal: walk away

and understand a consensus protocol

What is consensus?

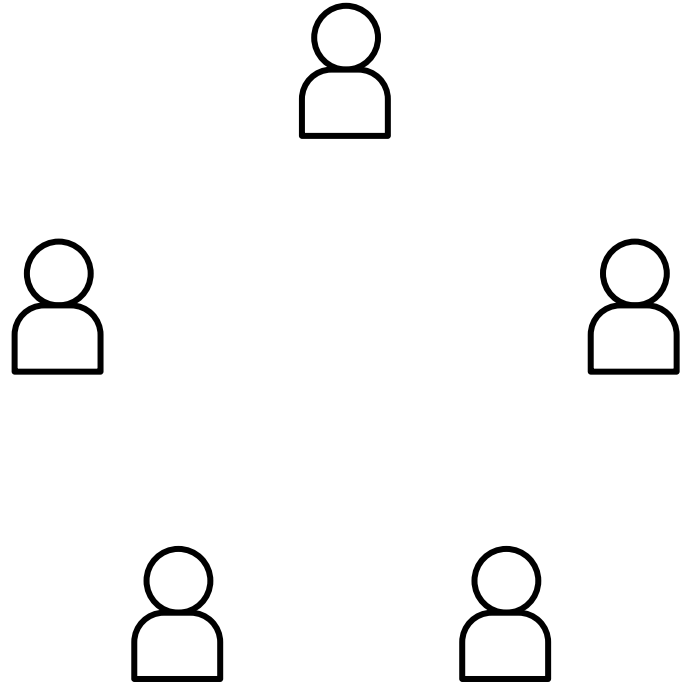
~~What is consensus?~~ Modeling Blockchain

Modeling Blockchain



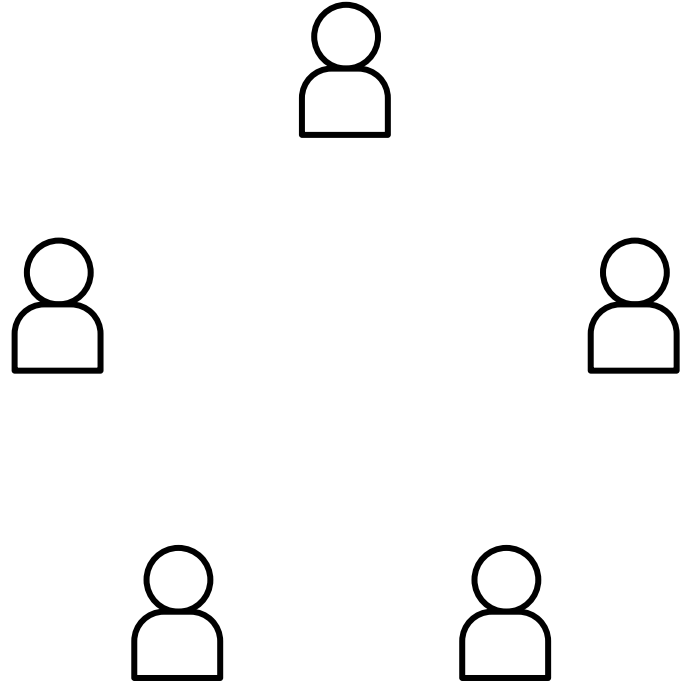
Modeling Blockchain

- Some *known* set of users
 - “permissioned”



Modeling Blockchain

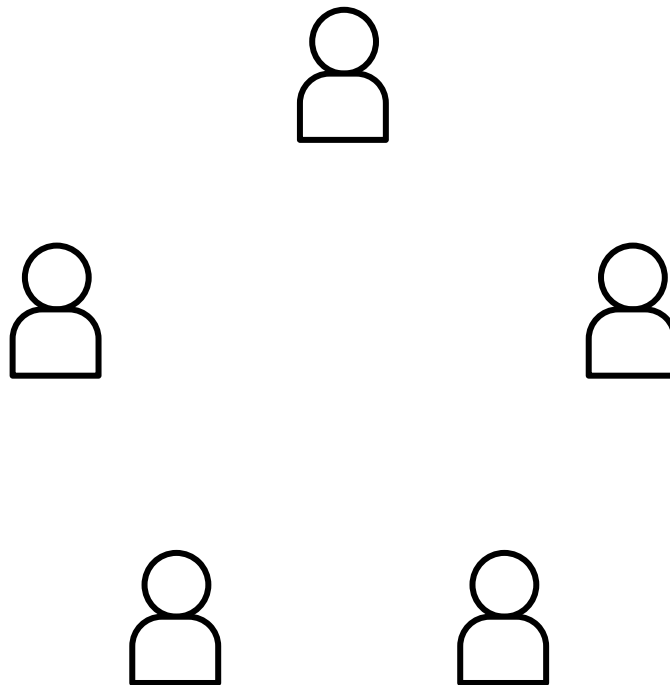
Why the permissioned setting?



Modeling Blockchain

Why the permissioned setting?

Answer: Proof-of-Stake

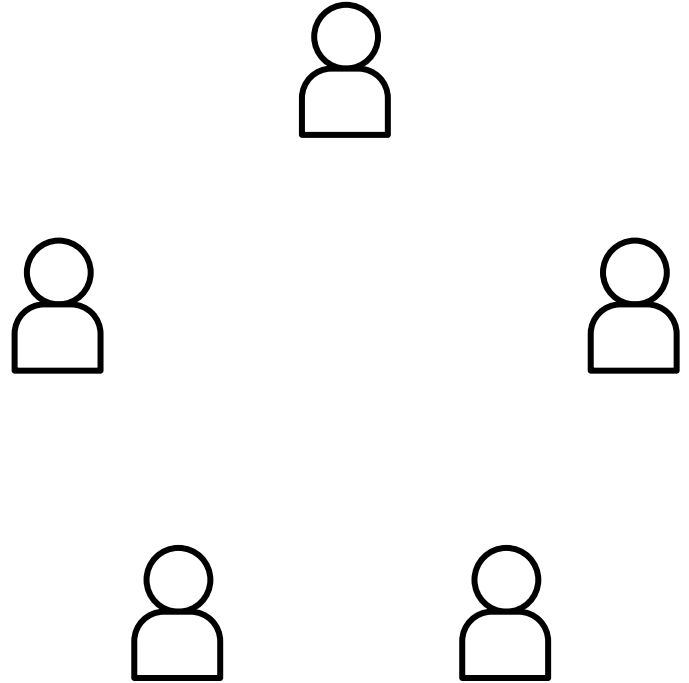


Modeling Blockchain

Why the permissioned setting?

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Different setting than
PoW!!

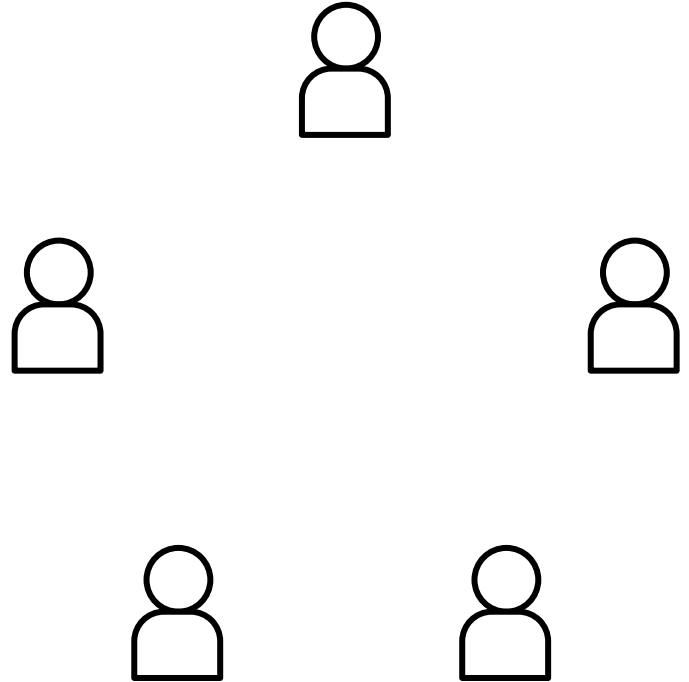


Modeling Blockchain

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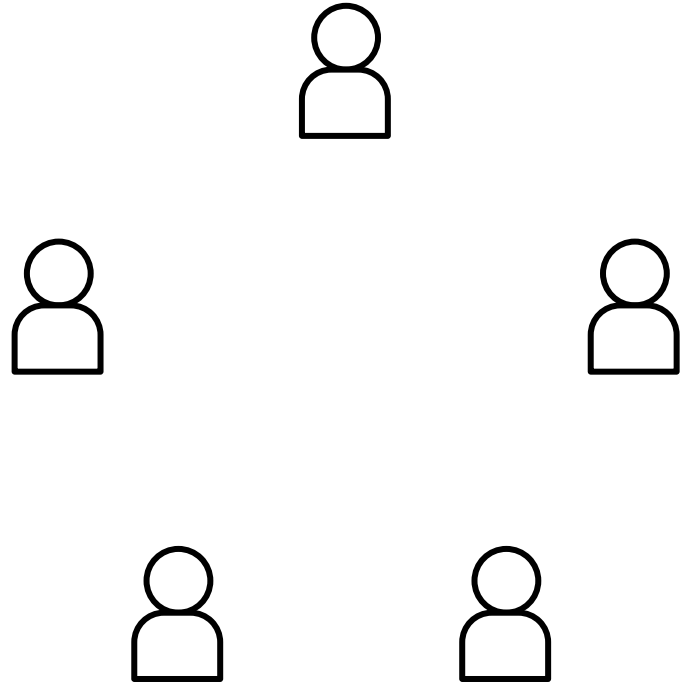
Answer: Proof-of-Stake

Different setting than PoW!! (true finality, speed, partition-resistant)



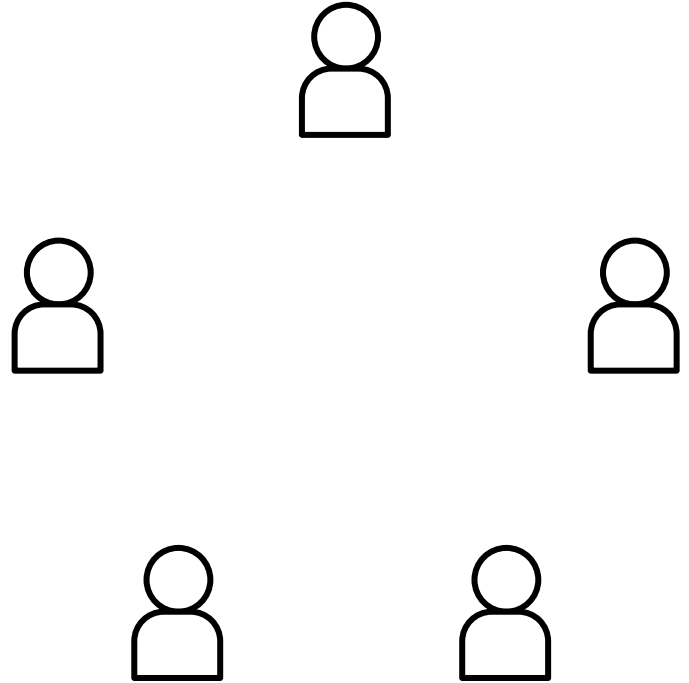
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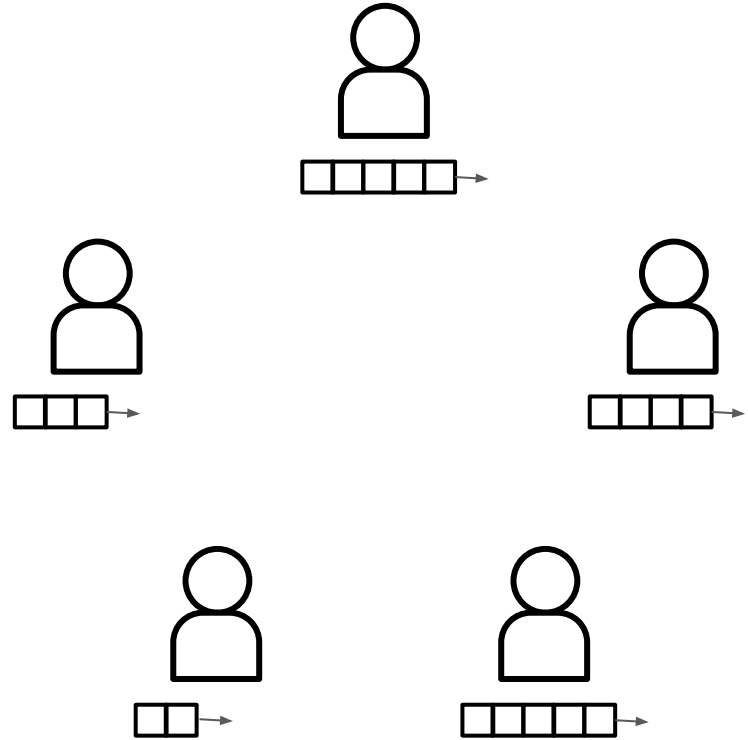
Modeling Blockchain

- Some *known* set of users
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- Each user maintains ordered chain of blocks



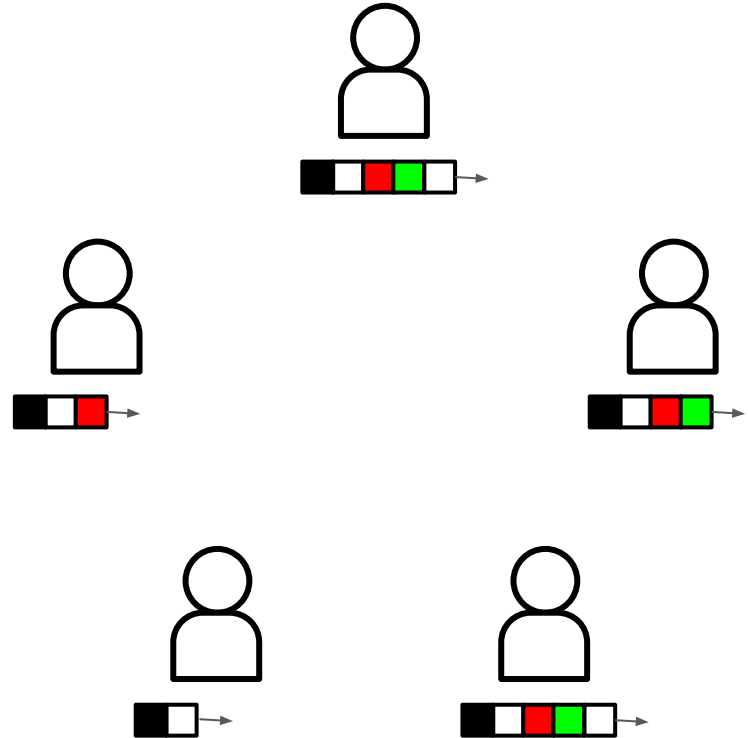
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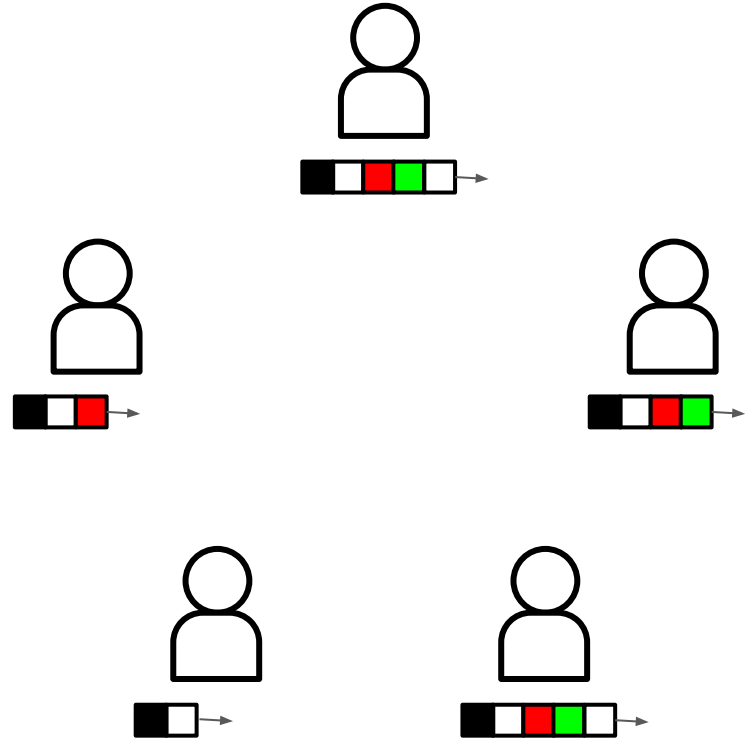
Modeling Blockchain

- Some *known* set of users
 - “permissioned”
- Each user maintains ordered chain of blocks
- Consistency: Everyone sees a prefix of the same chain!



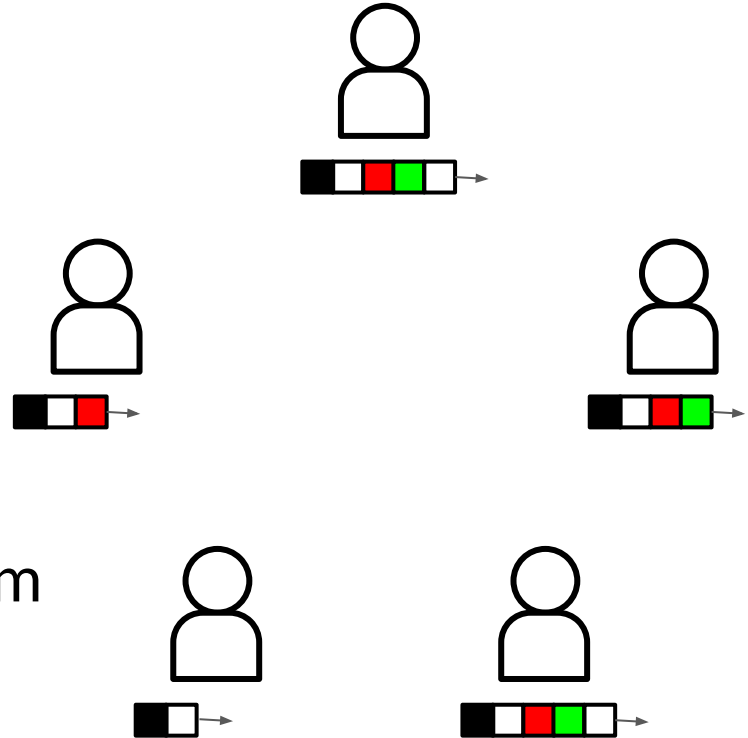
Modeling Blockchain

- Some *known* set of users
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- Each user maintains ordered chain of blocks
- Consistency
- Liveness

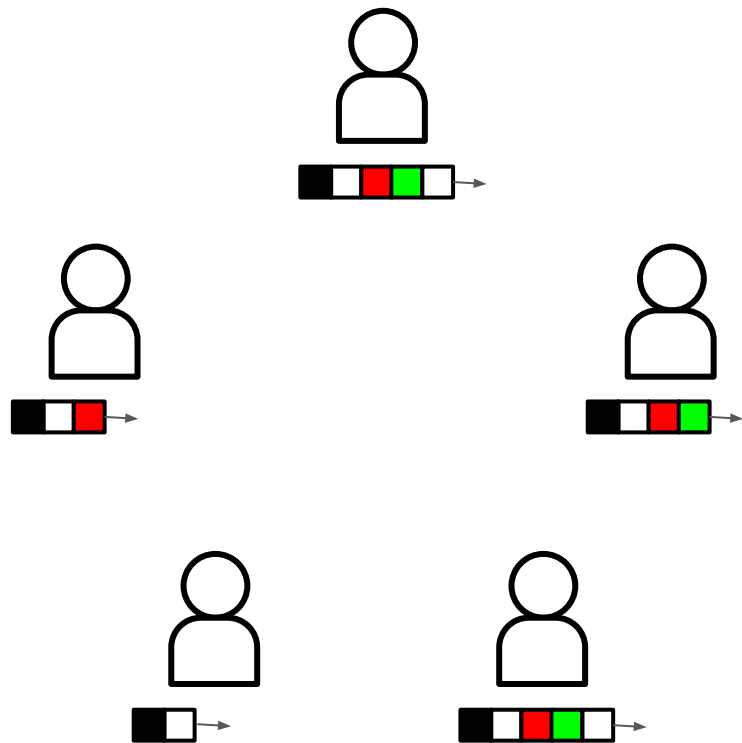


Modeling Blockchain

- Some *known* set of users
 - “permissioned”
- Each user maintains ordered chain of blocks
- Consistency
- Liveness: must be able to confirm new blocks

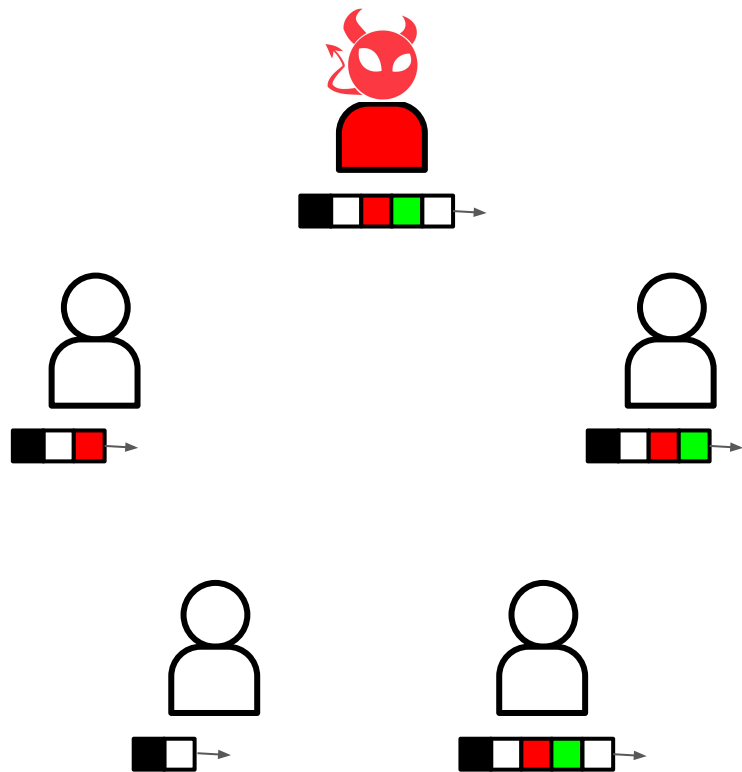


Introducing adversaries



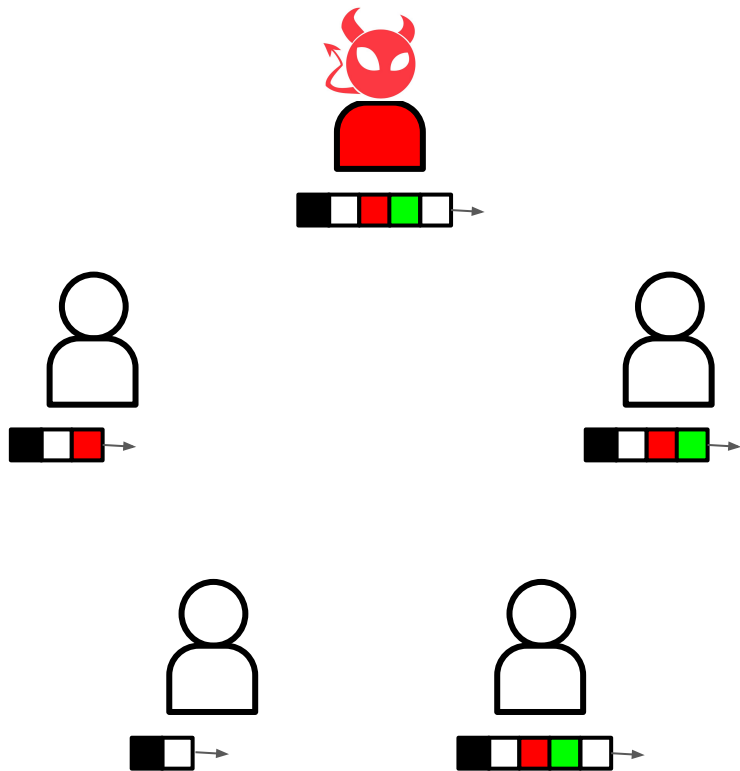
Introducing adversaries

- Malicious users



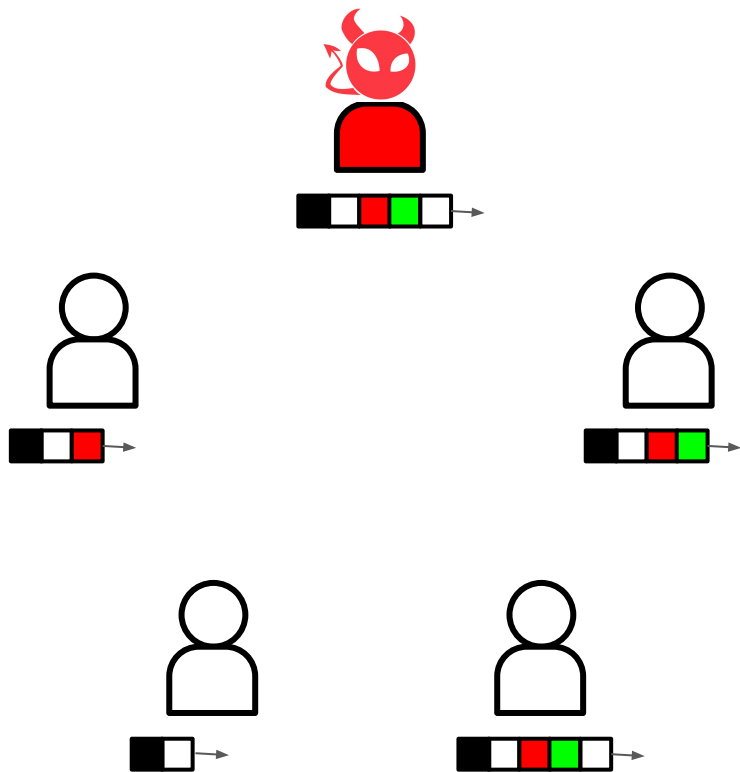
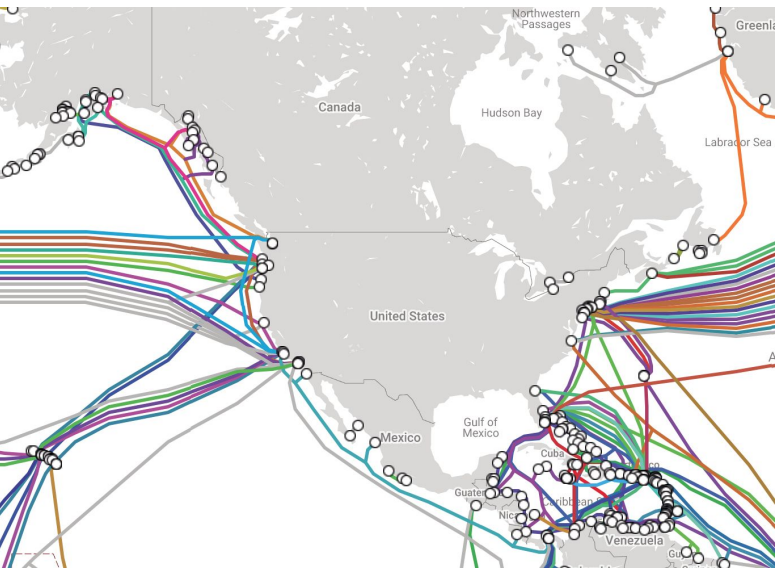
Introducing adversaries

- Malicious users
- Messages may be lost, delayed, reordered

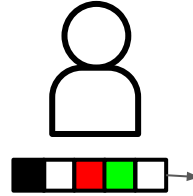
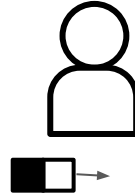
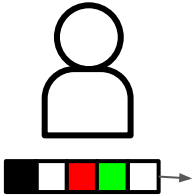
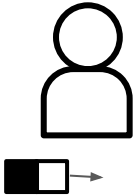
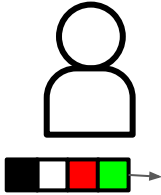
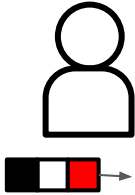
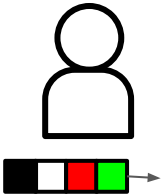
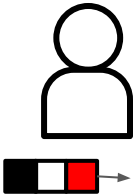
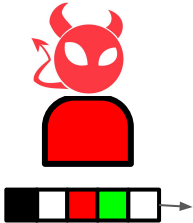
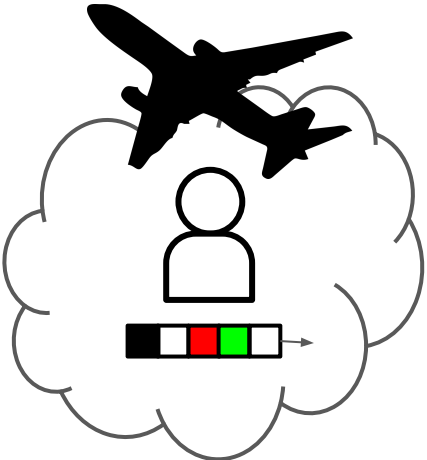


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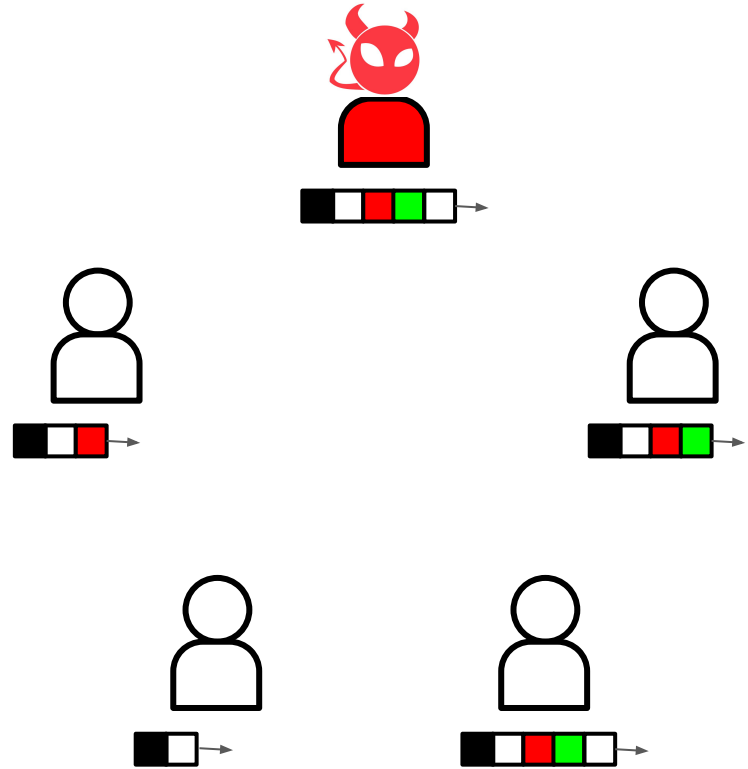
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???



This problem is notoriously hard!



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 - Casper
 - Algorand
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**New “streamlined”
paradigms...**

**...but can we eliminate
the subtleties?**

Motivating Simpler Consensus Protocols

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- Simpler Implementation

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- Fewer Bugs

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Motivating Simpler Consensus Frameworks

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Motivating Simpler Consensus F

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Our Work: Streamlet

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Goal:

A “Simplest Possible”,
Easy-to-Understand,
Textbook Consensus Protocol

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Goal:

A “Simplest Possible”,
Easy-to-Understand,
Textbook Consensus Protocol (Blockchain)

Our Work: Streamlet

Two Assumptions:

1. **Epochs**

Processes have local clocks,
and run in synchronized* epochs of 1 sec each.

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Two Assumptions:

1. **Epochs**

Processes have local clocks,
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2. **Elect a leader in each epoch, known by all**

i.e. randomly chosen, given epoch e

$$L_e = H(e) \bmod n$$

Assumptions:

- ❑ (Synchronized*)
epochs of length 1 sec
- ❑ Each epoch
has random leader

Definitions

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pointer to parent block

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Epoch number in which
block was 'proposed'

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Assumptions:

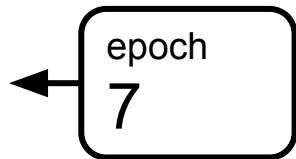
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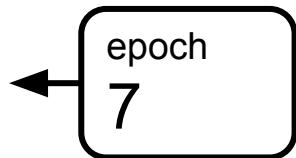


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- Notarized block

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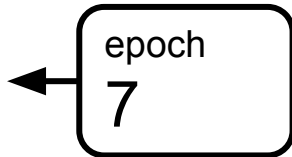


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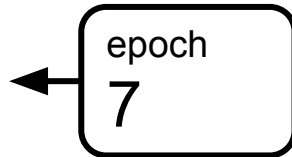


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(implies a majority of honest processes have signed it)

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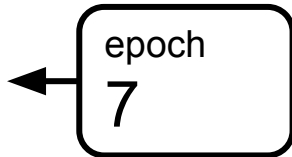


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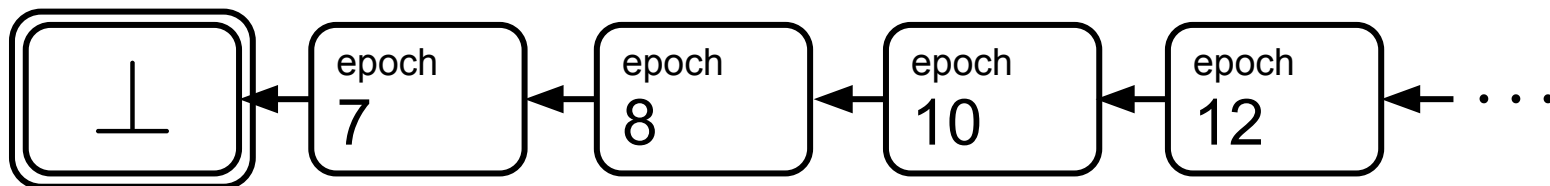


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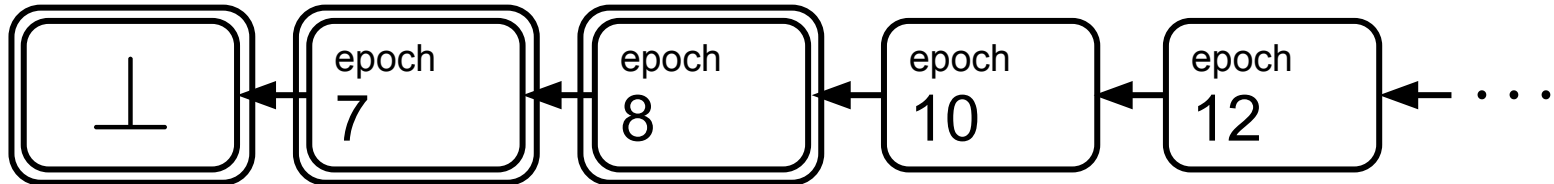


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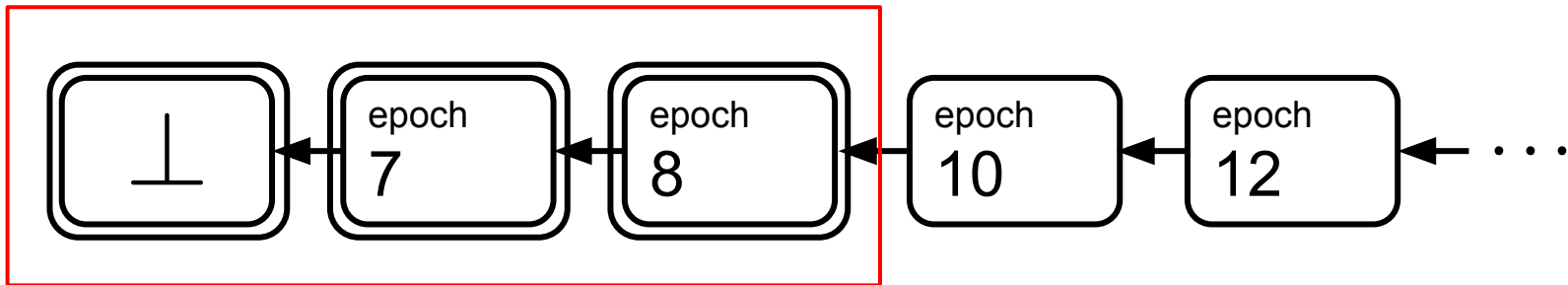


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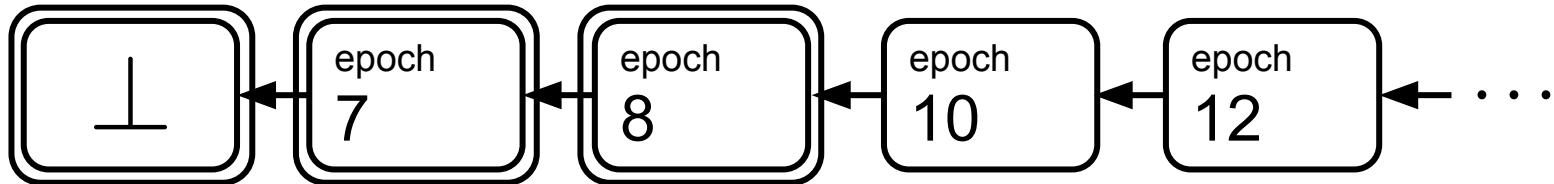


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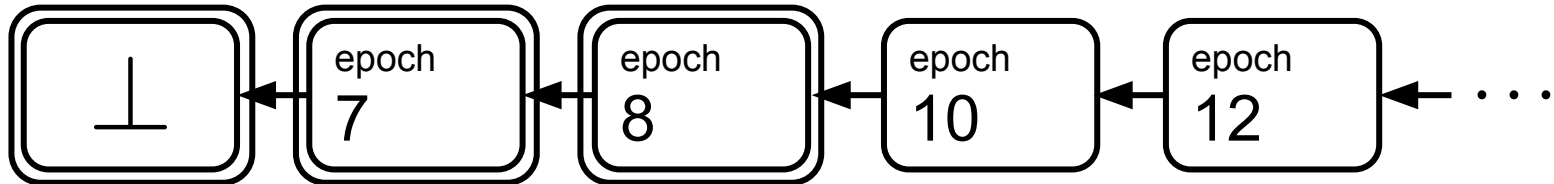


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- Block "height" \neq epoch #

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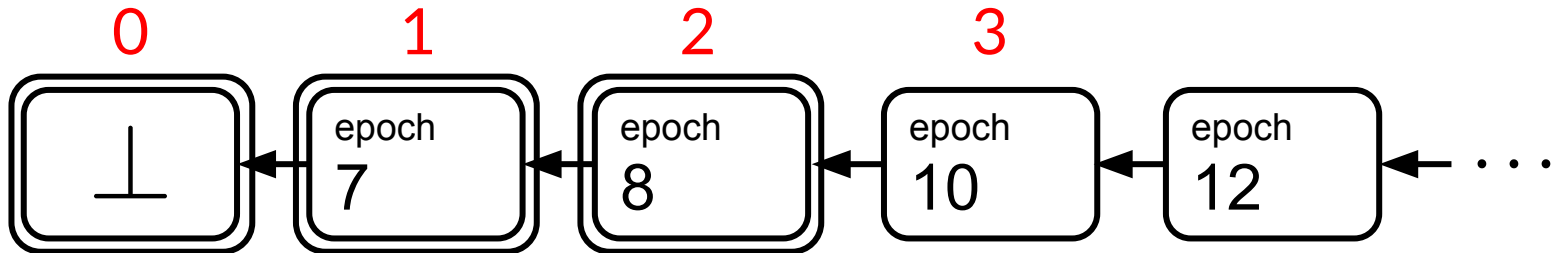


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The Streamlet Protocol

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In every epoch $e = 1, 2, \dots$

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In every epoch $e = 1, 2, \dots$

- leader,
creates a new block $b = (H(b'), e, \text{txs})$
extending longest notarized chain they've seen so far

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- voters,
signs first proposal b (from leader, for e)
i.f.f. b extends a longest notarized chain seen so far (by voter)

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finalization rule:

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take any notarized chain
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finalization rule:

take any notarized chain
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chop off the last block,
and finalize

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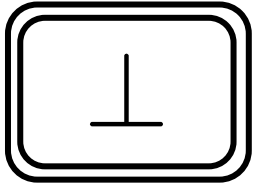
Example

In every epoch $e = 1, 2, \dots$

- ❑ leader proposes $b = (H(b'), e, \text{txs})$ extending longest notarized chain they've seen
- ❑ voters sign the first valid proposal b ,
but i.f.f. b also extends a longest notarized chain the *voter* has seen (notarized=2/3 votes)

finalize any notarized chain ending with 3 consecutive epochs, chopping off last block

Example

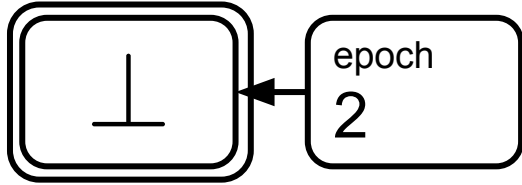


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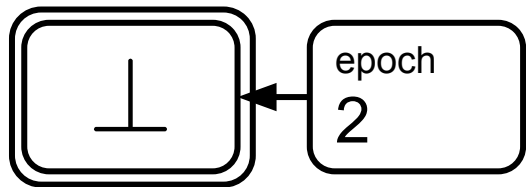


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Example (3 honest, 1 malicious)

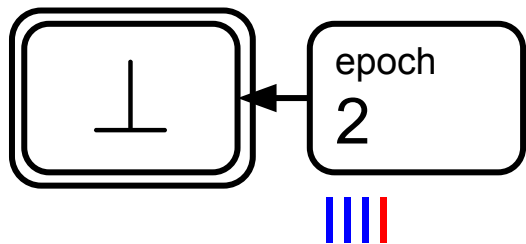


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finalize any notarized chain ending with 3 consecutive epochs, chopping off last block

Example (3 honest, 1 malicious)

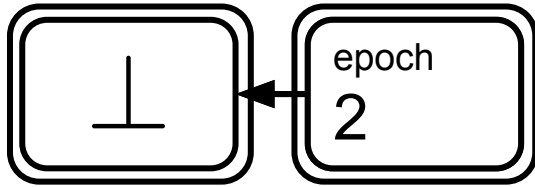


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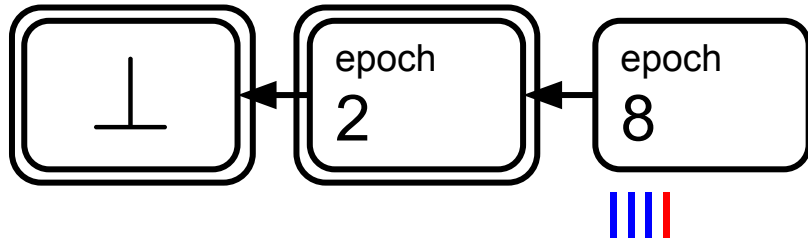


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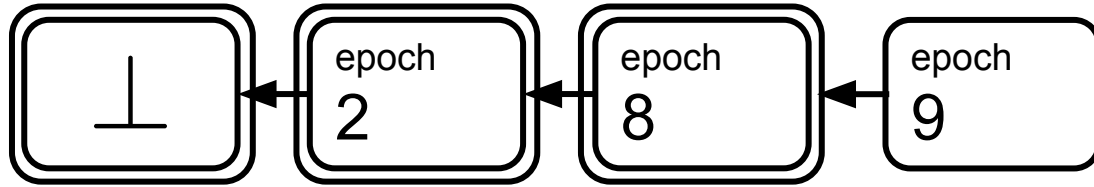


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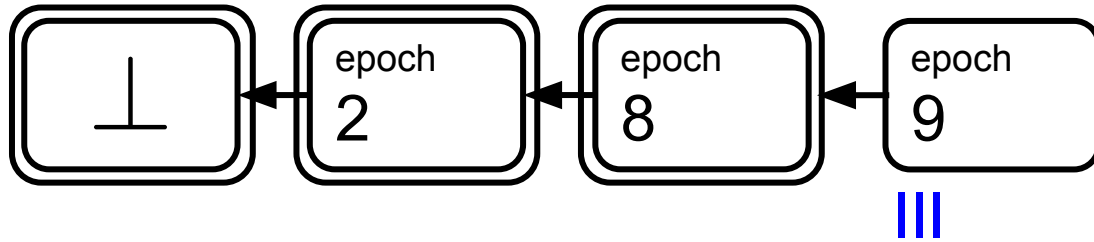


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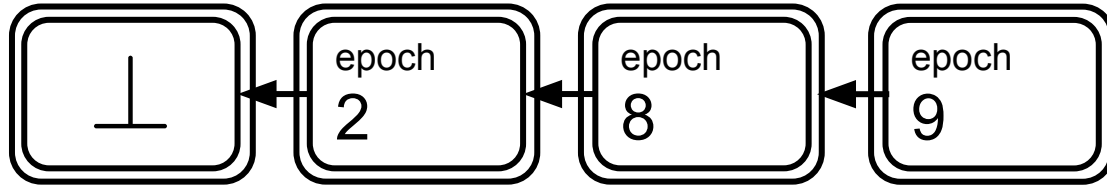


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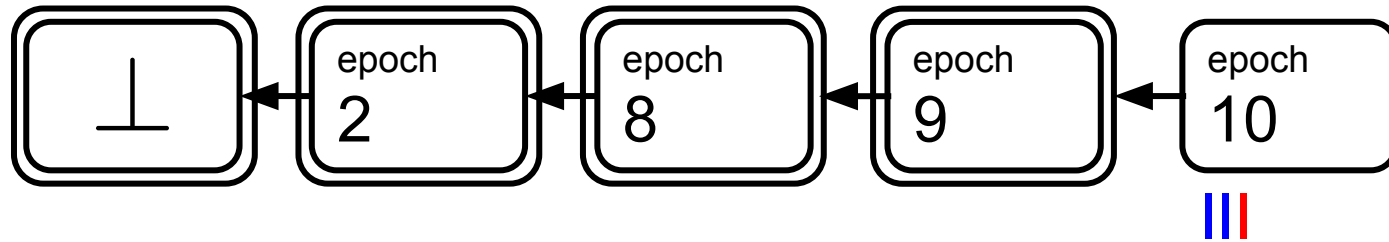


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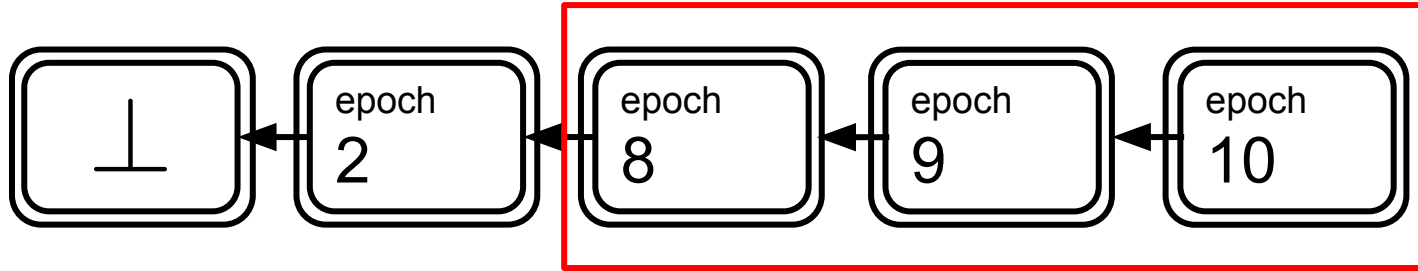


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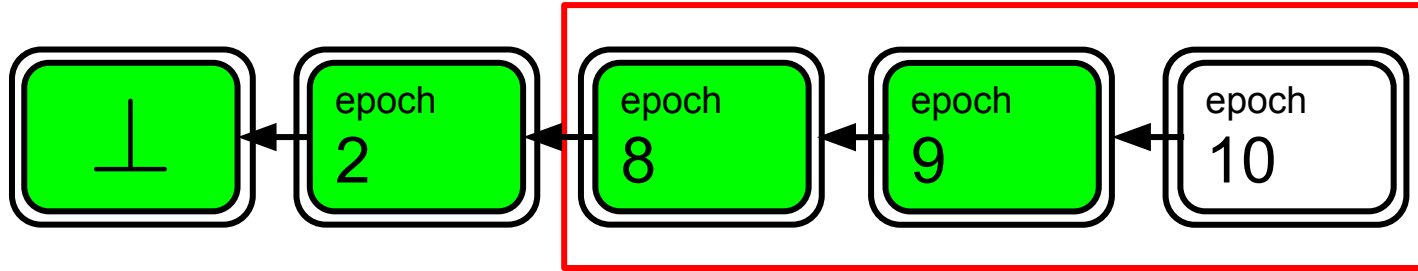


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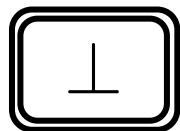
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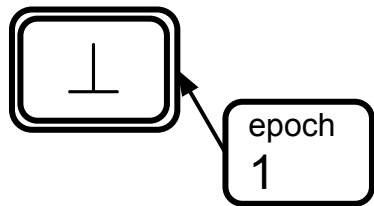


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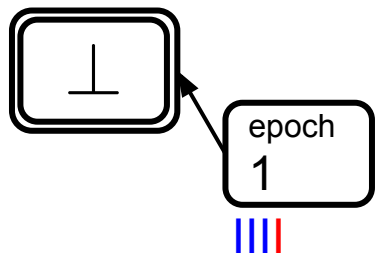


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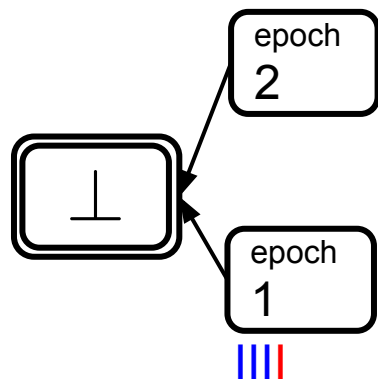


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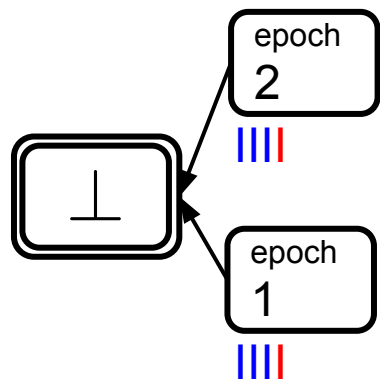


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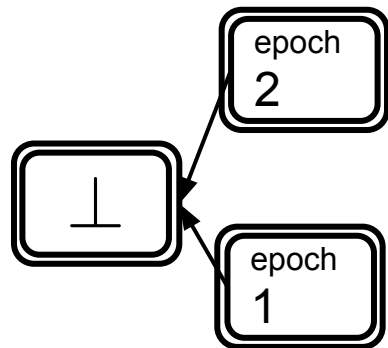


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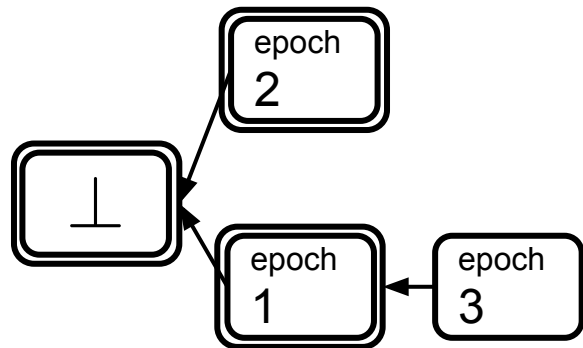


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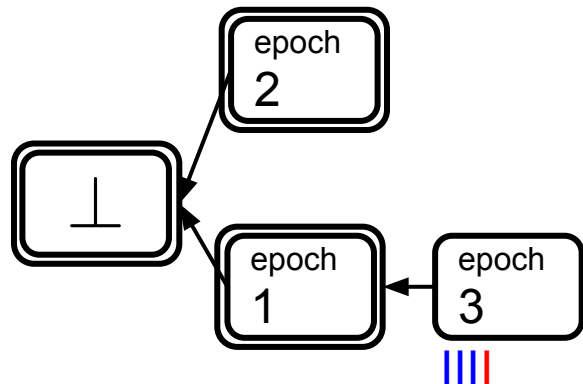


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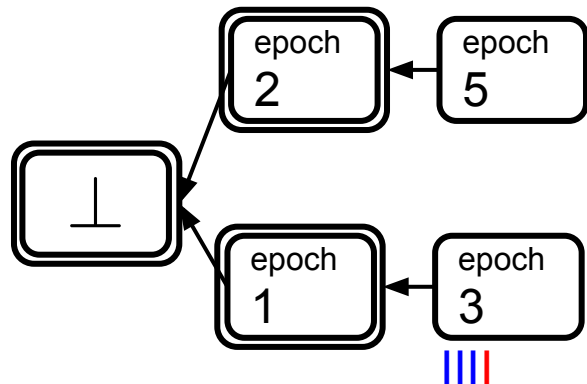


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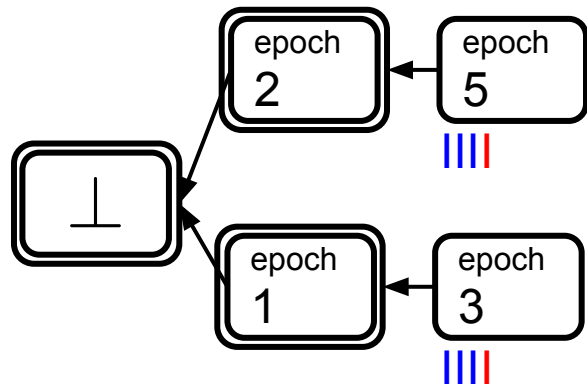


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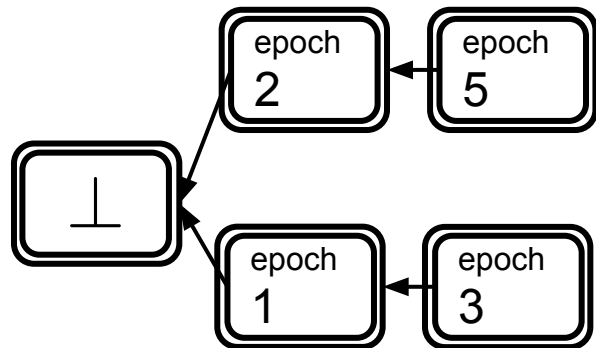


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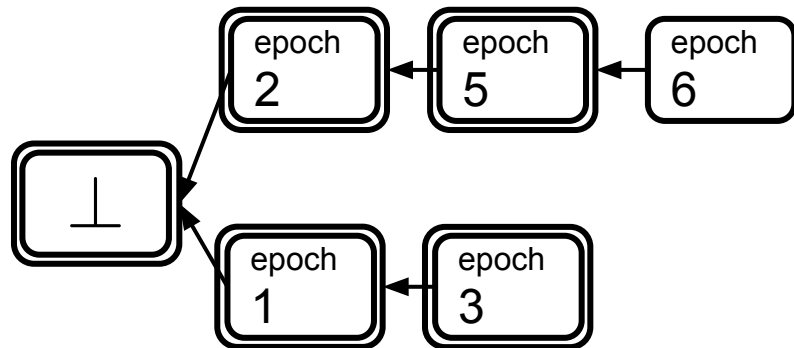


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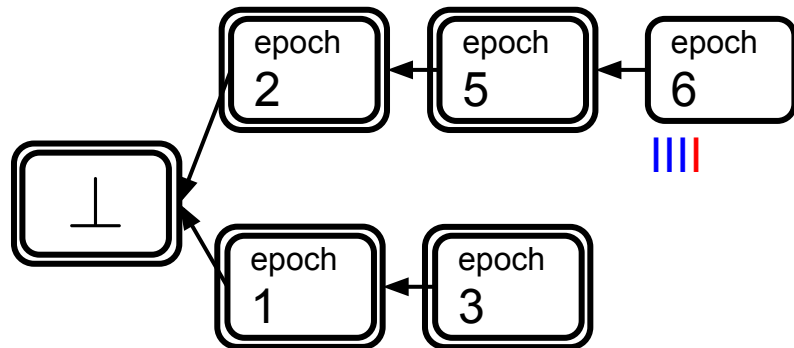


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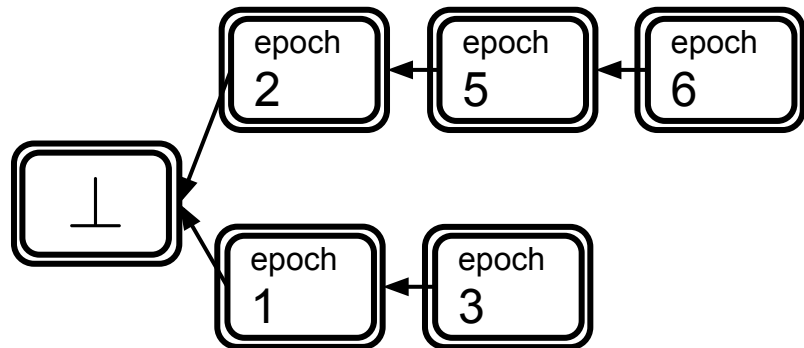


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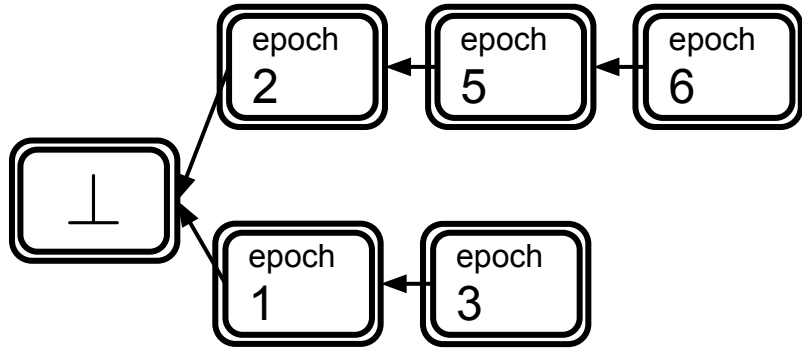


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Allows two notarized blocks at the same height!

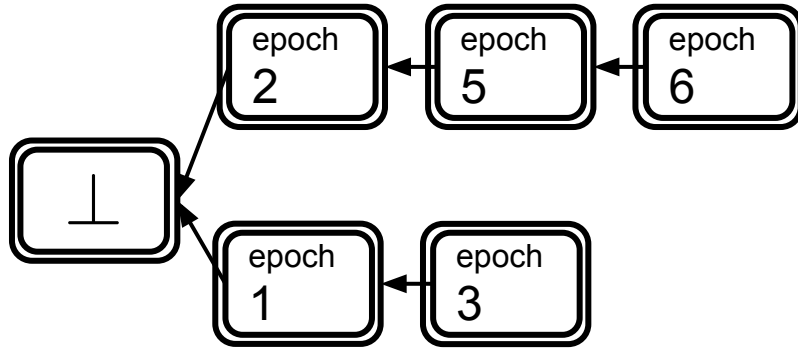


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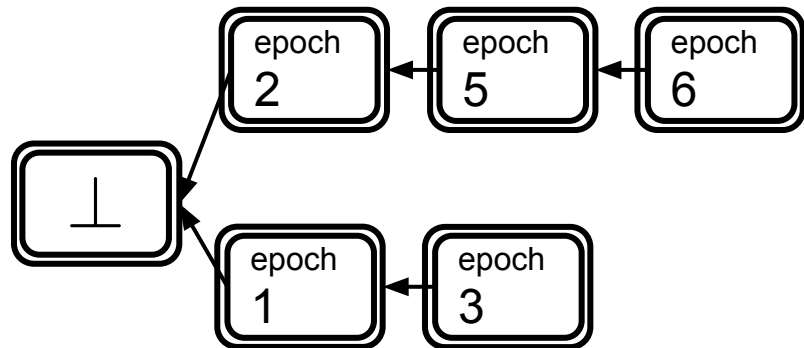


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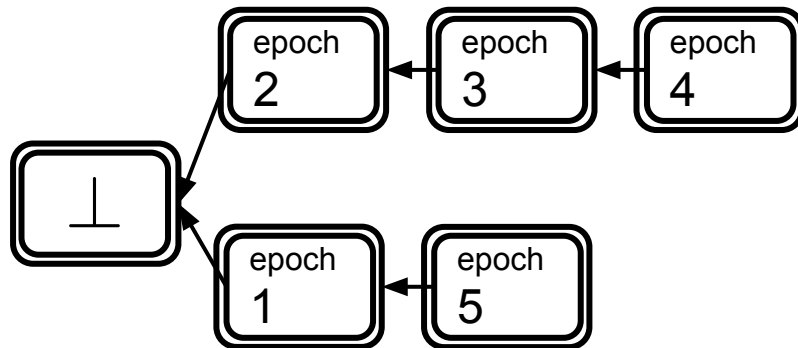
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Impossible

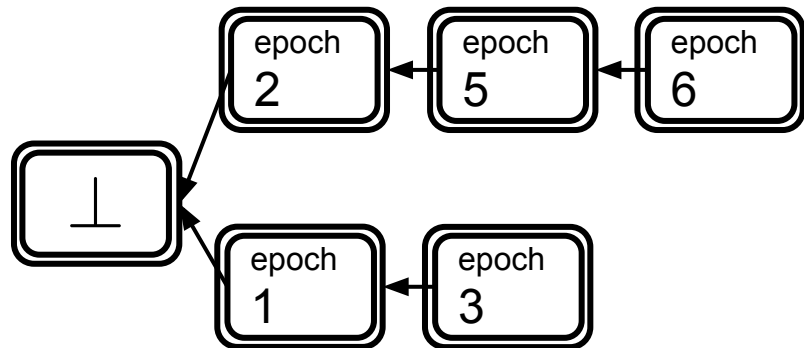


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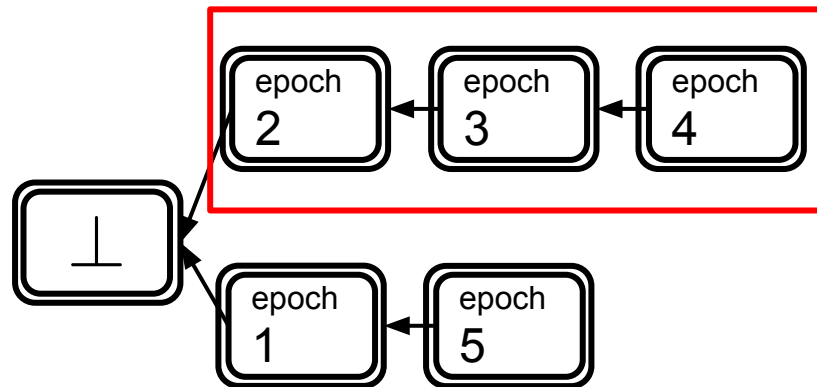
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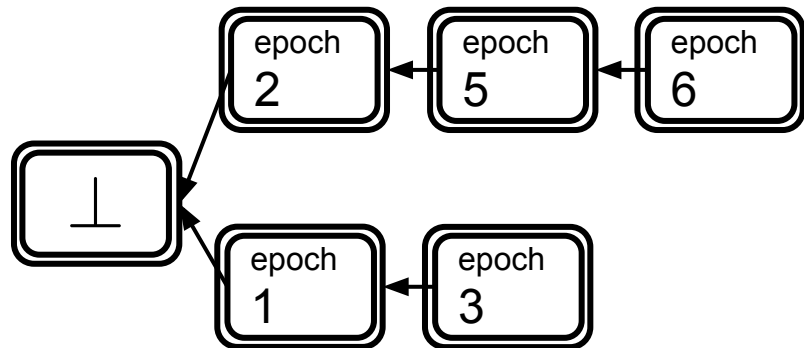


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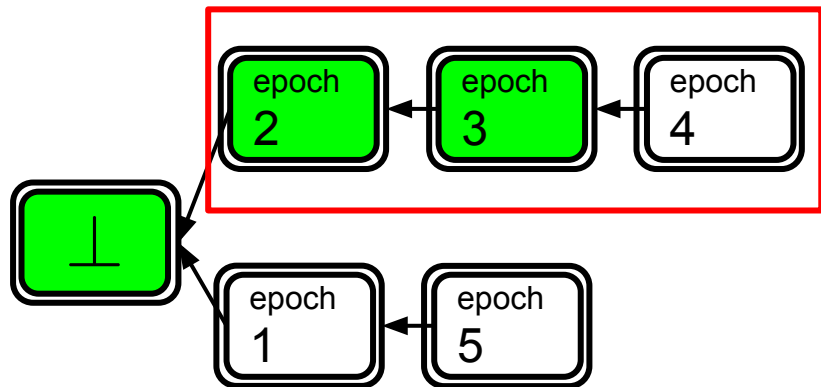
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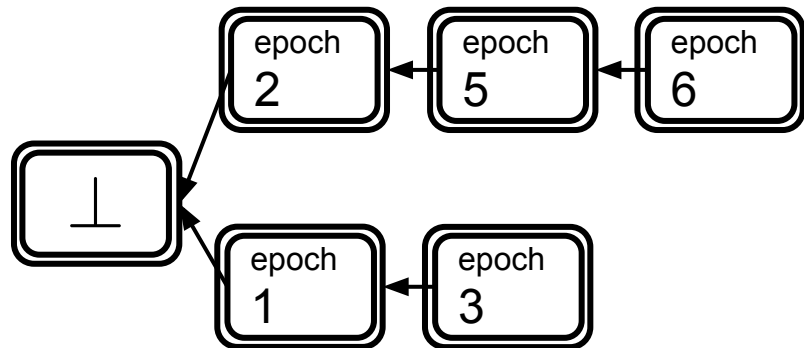


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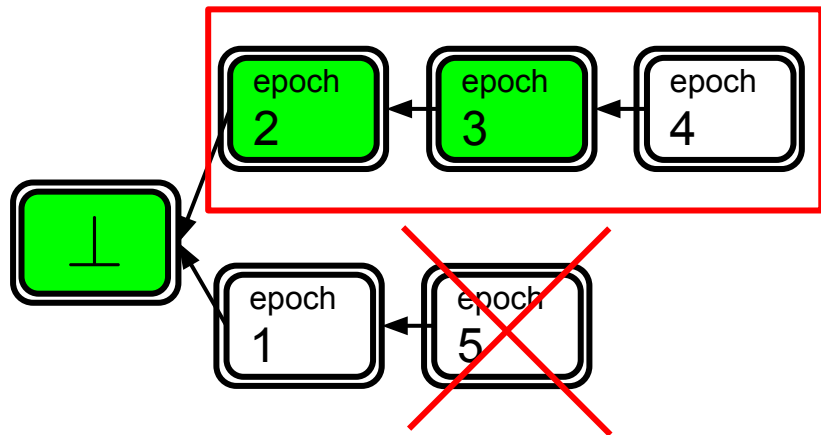
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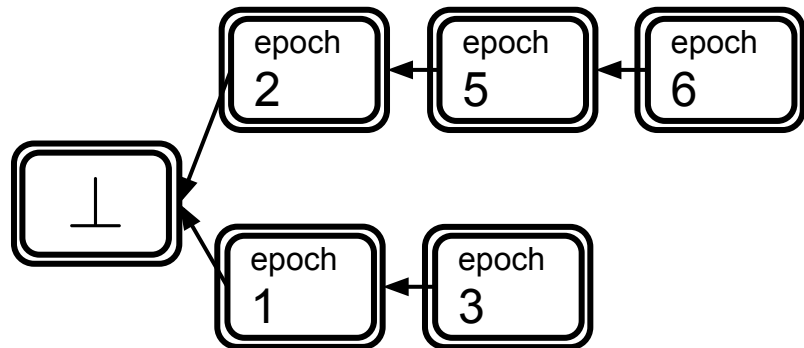


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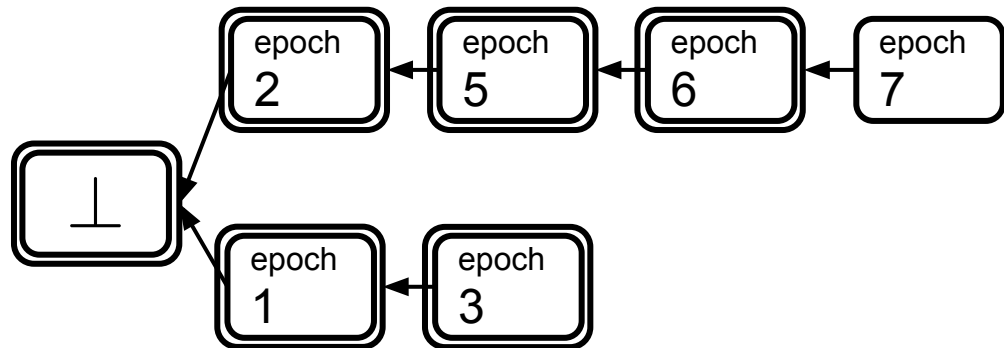


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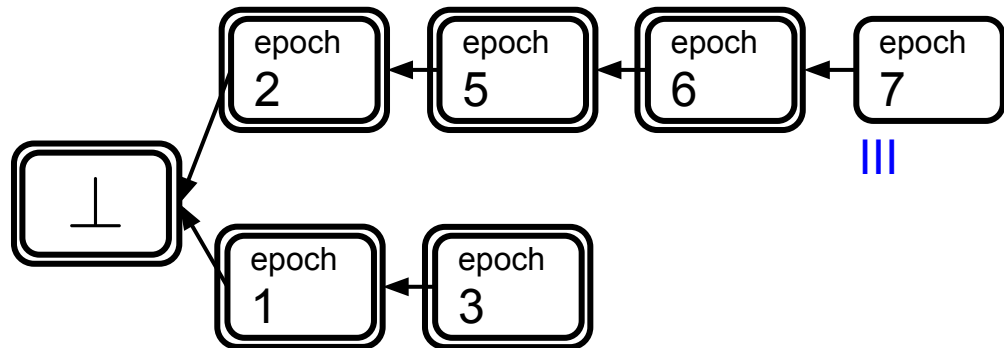


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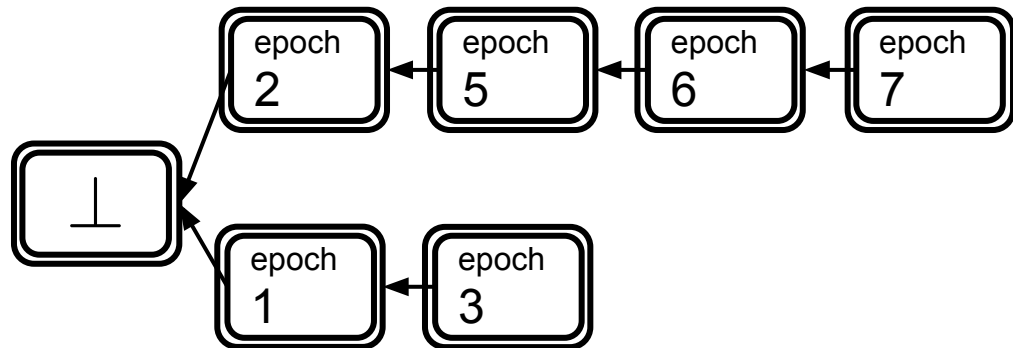


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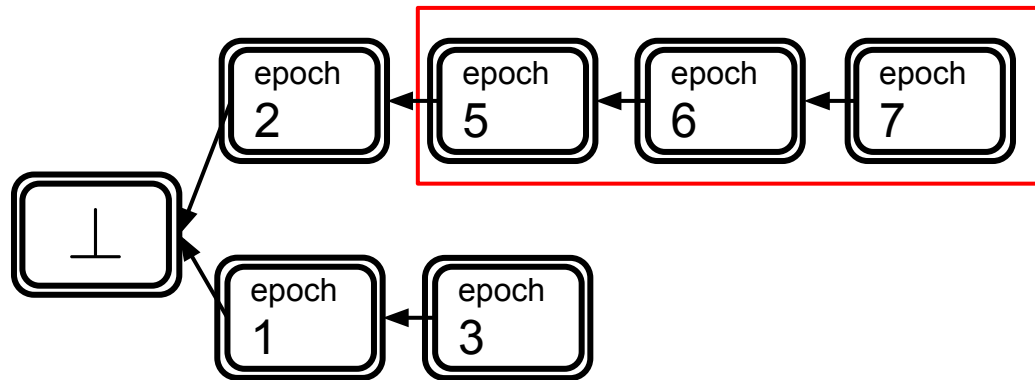


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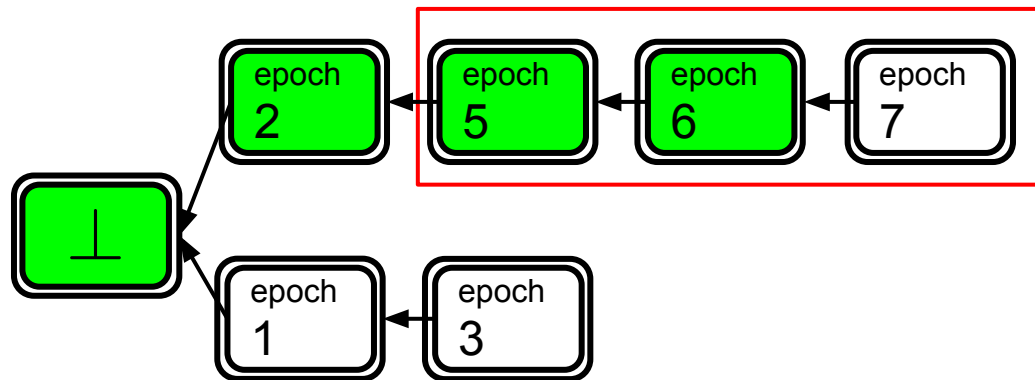


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finalize any notarized chain ending with 3 consecutive epochs, chopping off last block

Consistency Sketch

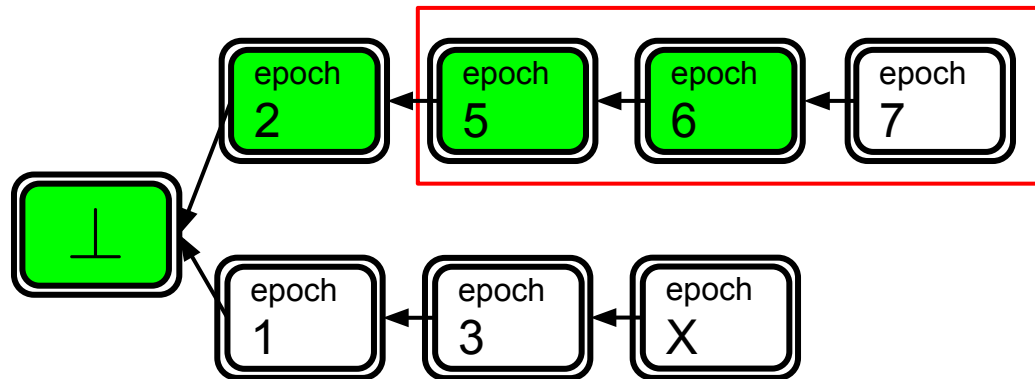


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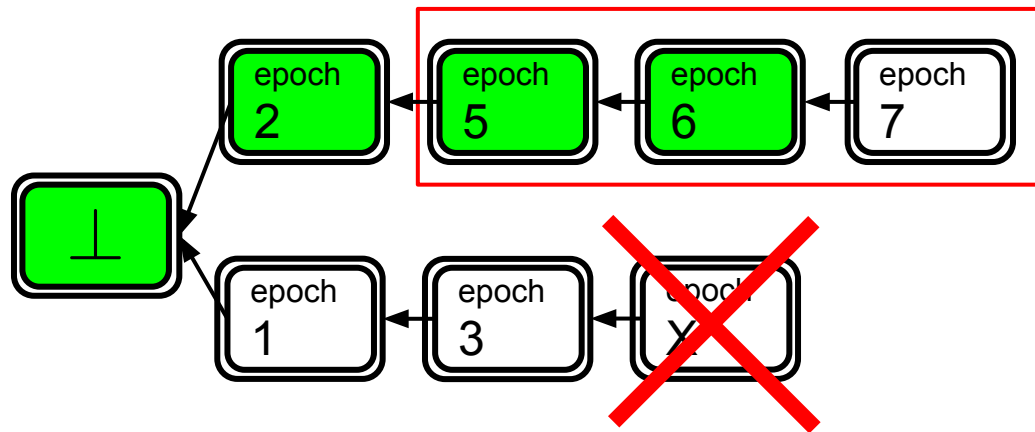


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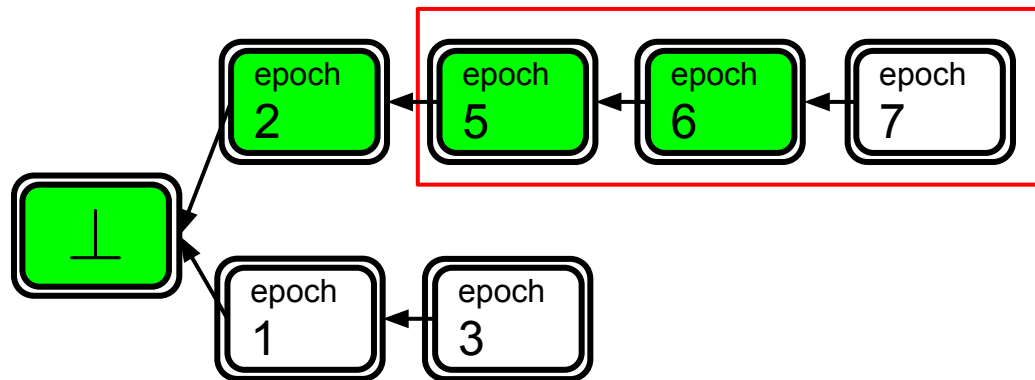


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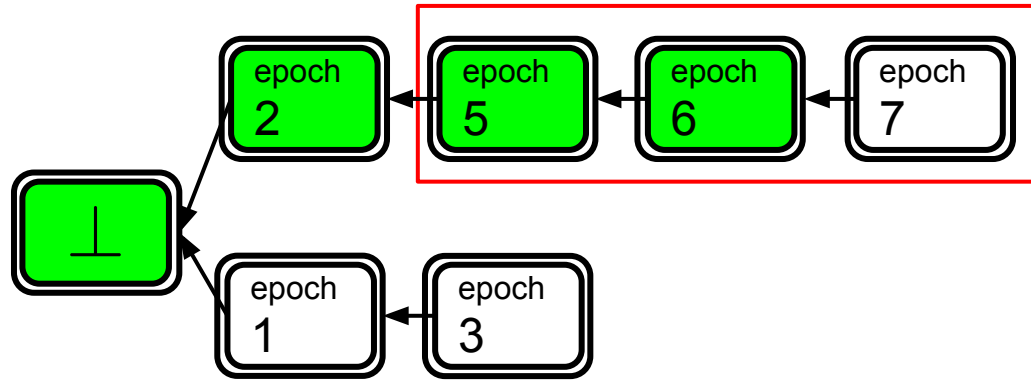


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Lemma 1: *Each epoch is associated with at most one notarized block.*

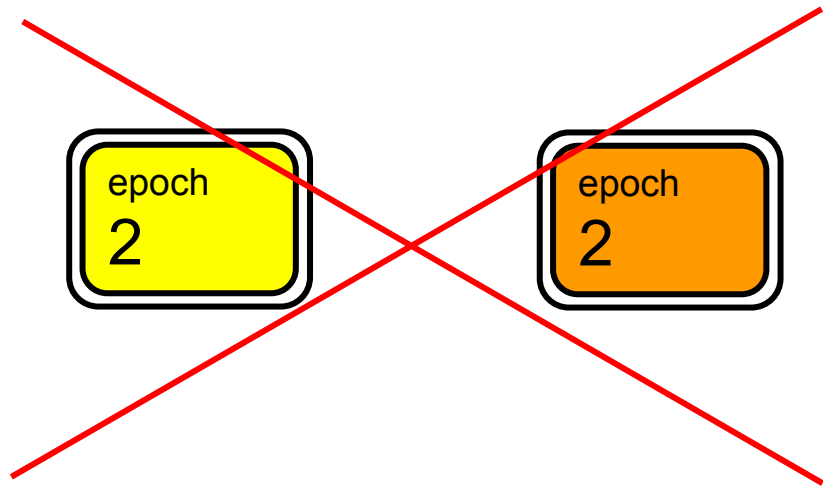
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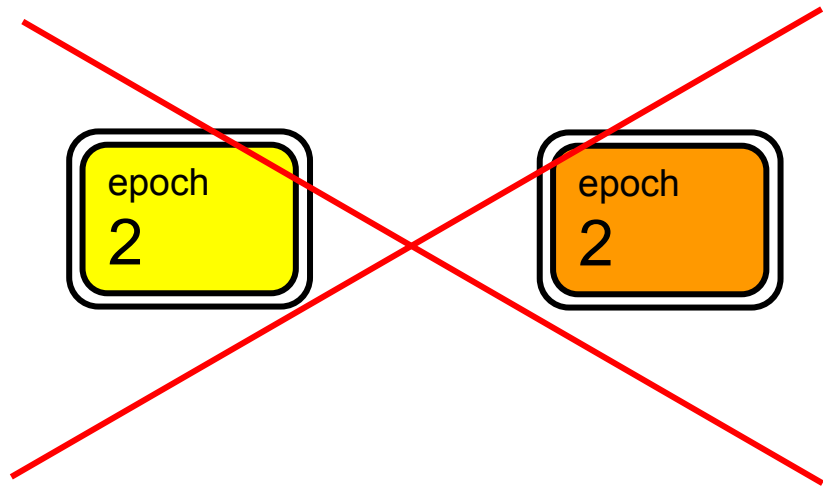
Cannot both be notarized!!



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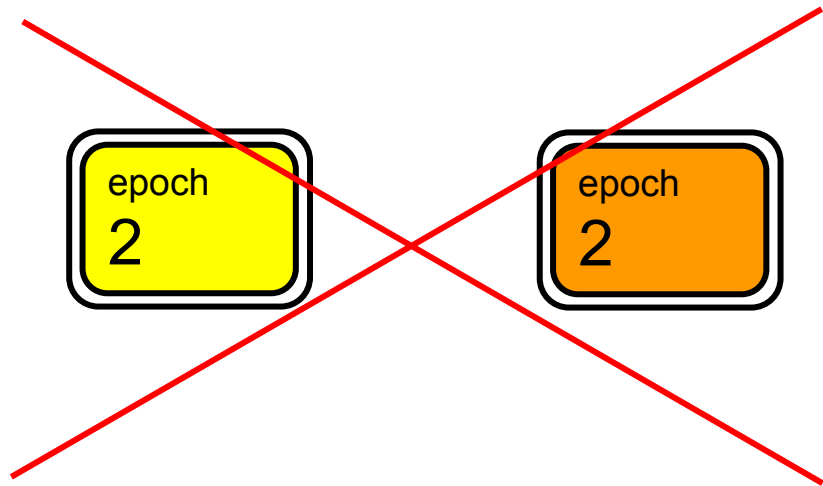
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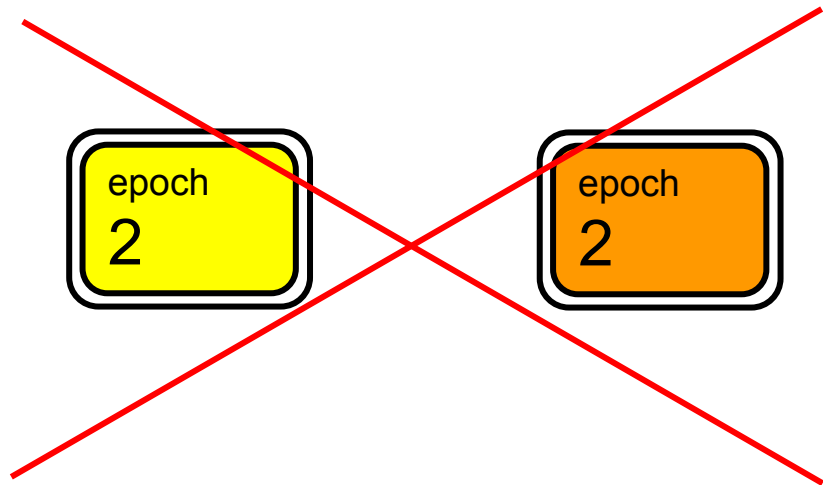
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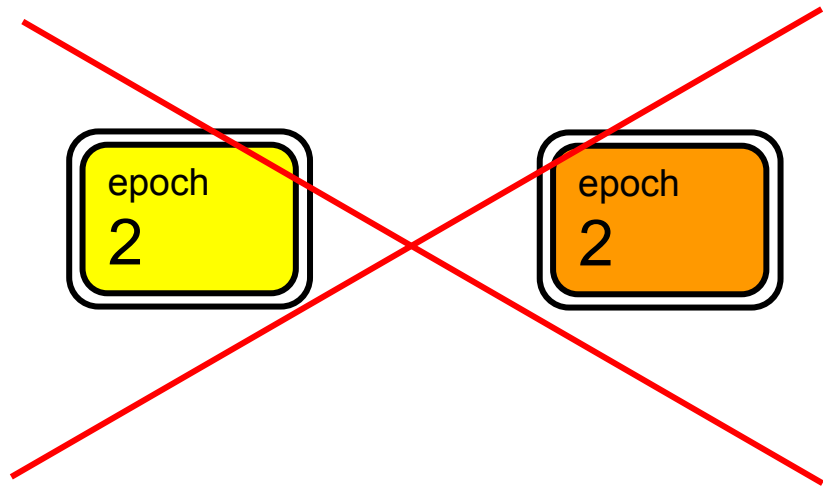
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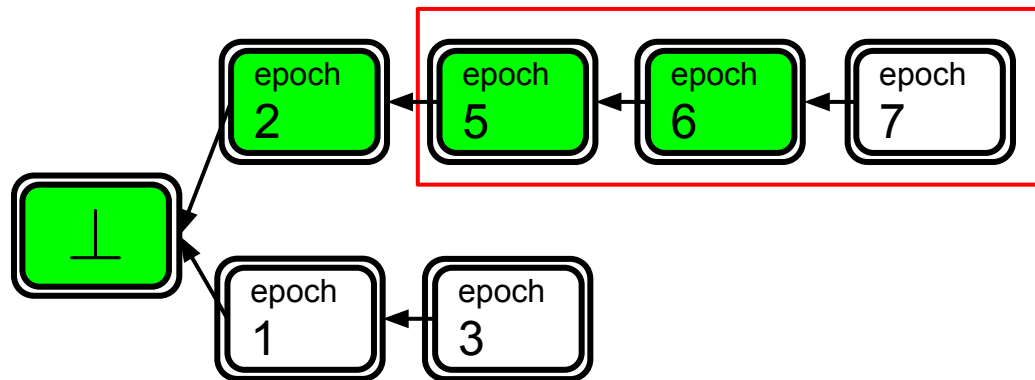


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- ❑ Letting $f < n/3$, we have (at best) $2n/3 + 2f < 4n/3$ votes to go around.

Consistency Sketch



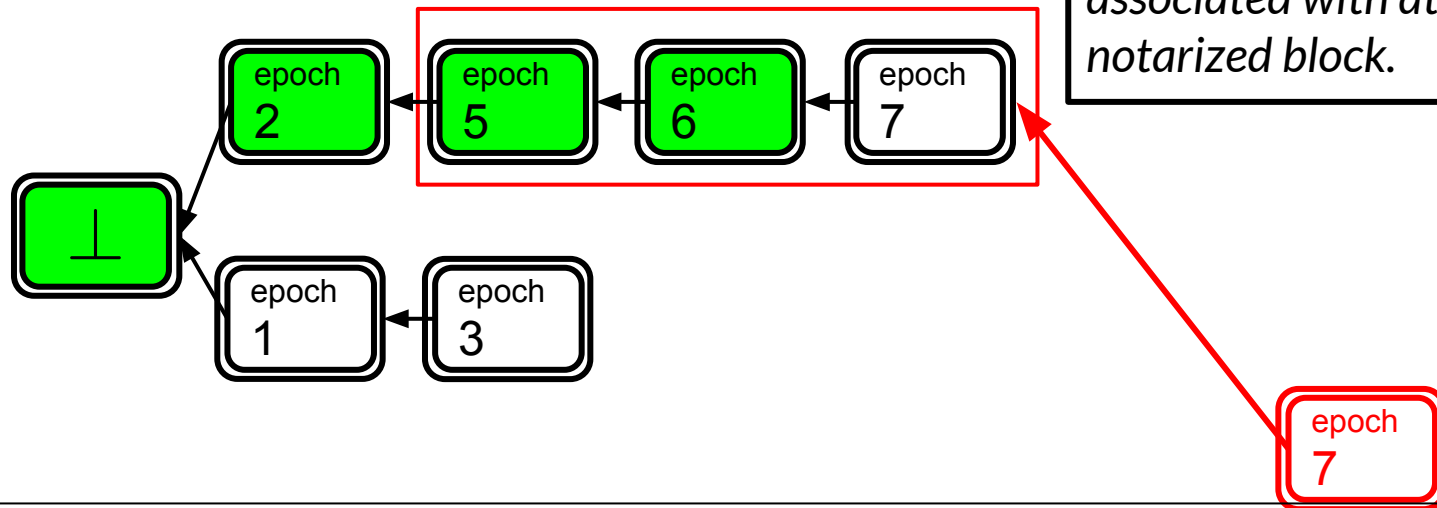
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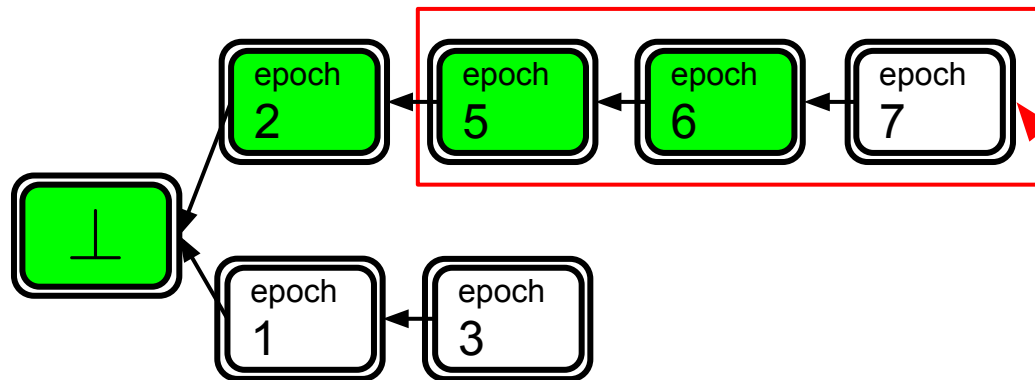


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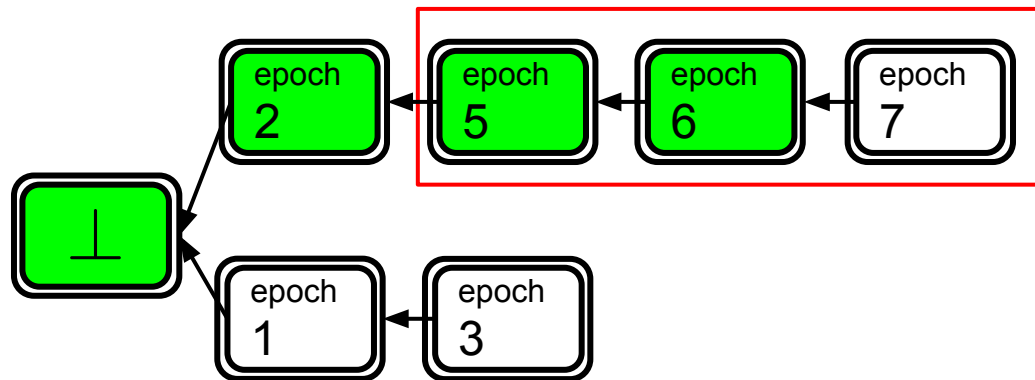


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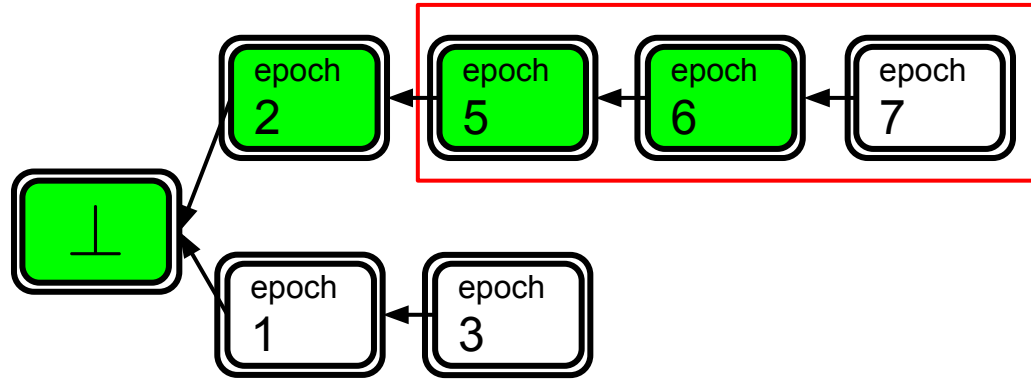


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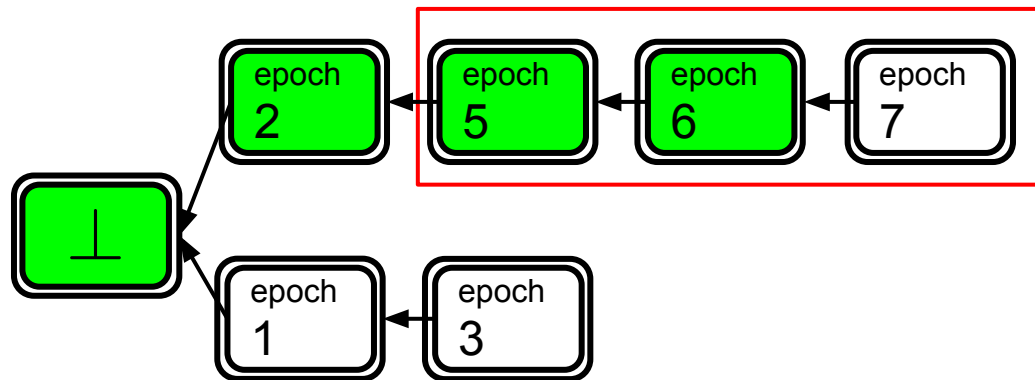
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Consistency Sketch



Now, the main lemma...

Consistency Sketch

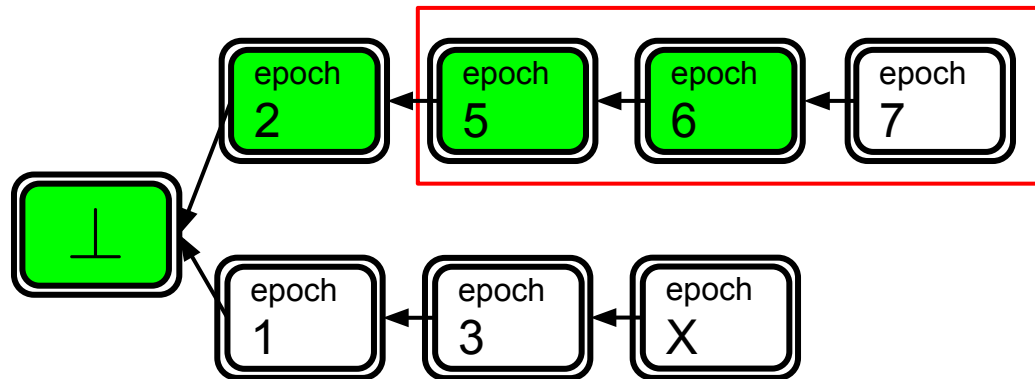


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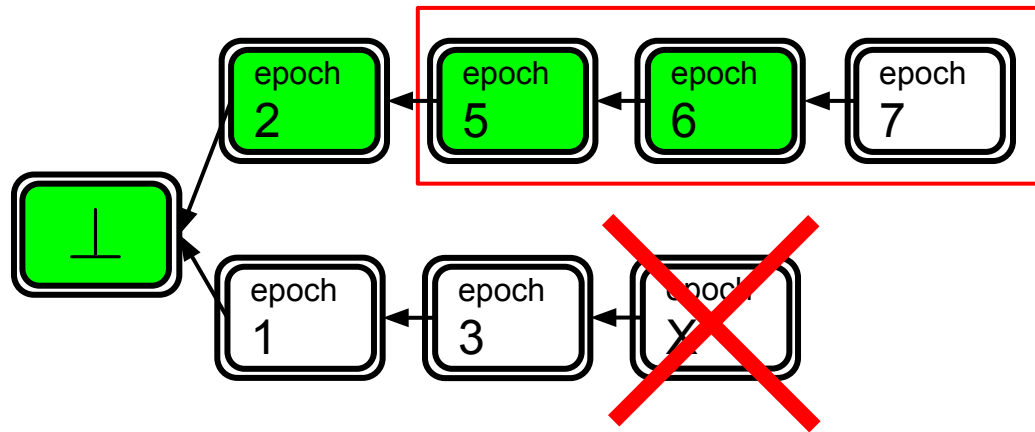


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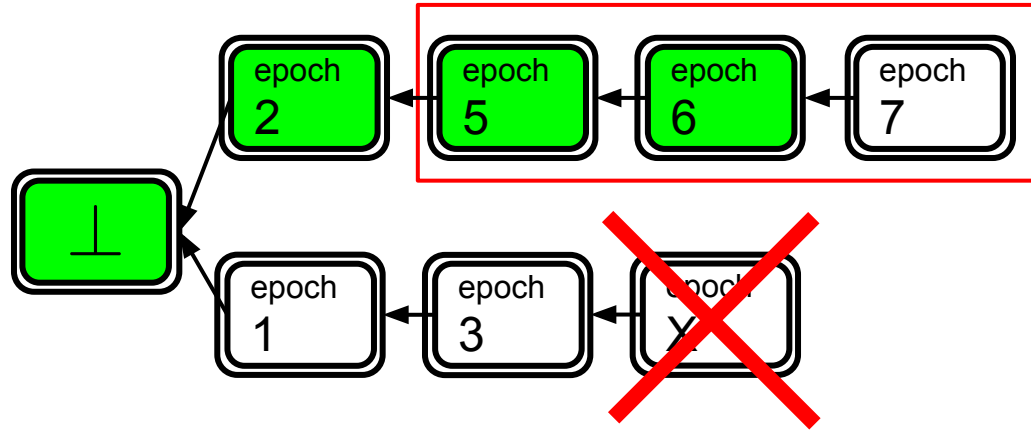


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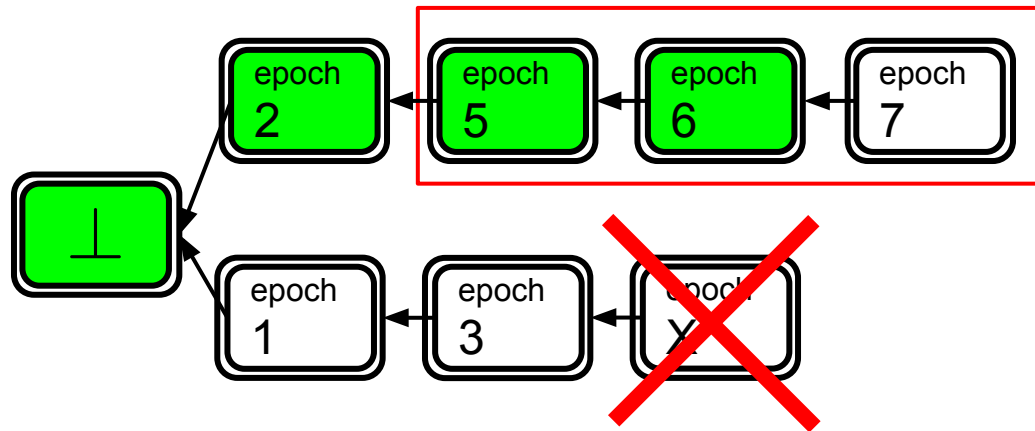
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Consistency Sketch



Lemma 2: No other notarized block, in past or future, can share the same height as block 6

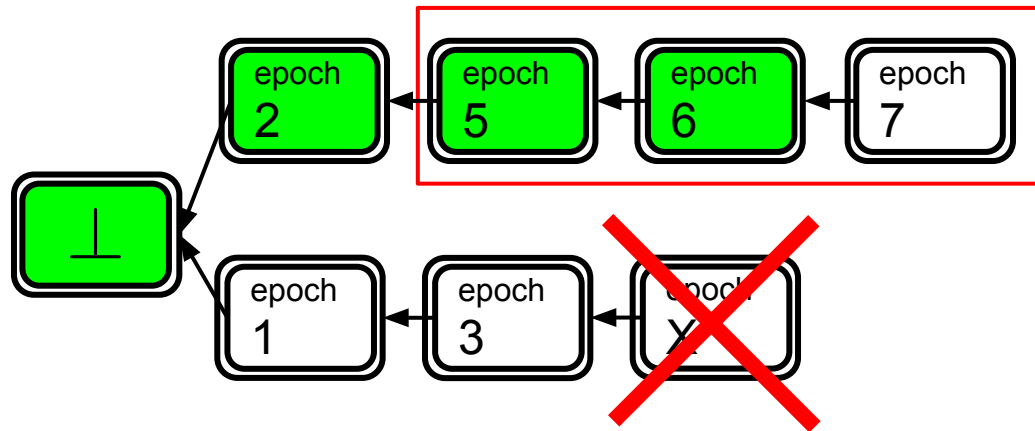
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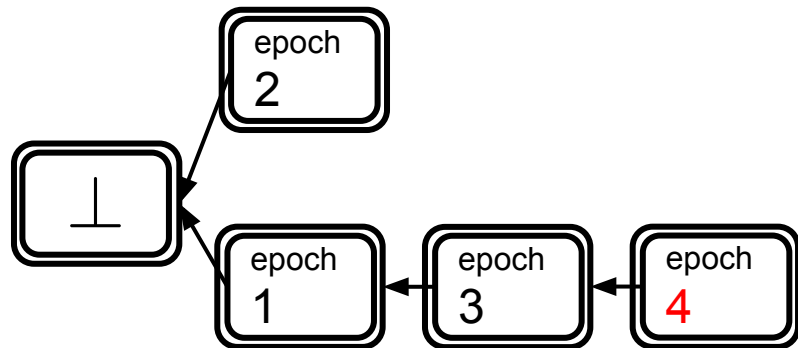
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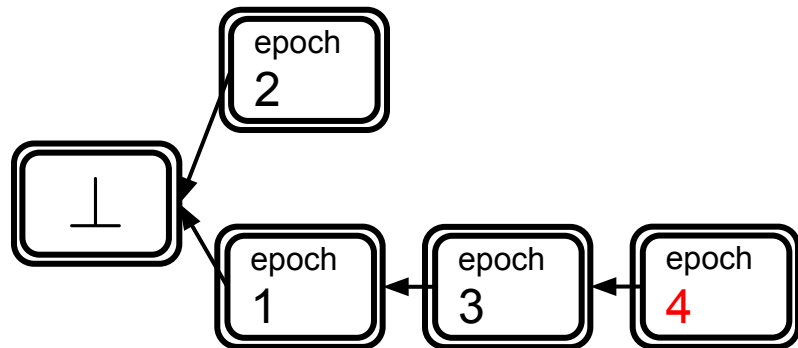
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4 processes, 3 honest, 1 malicious
Require 3 votes to notarize

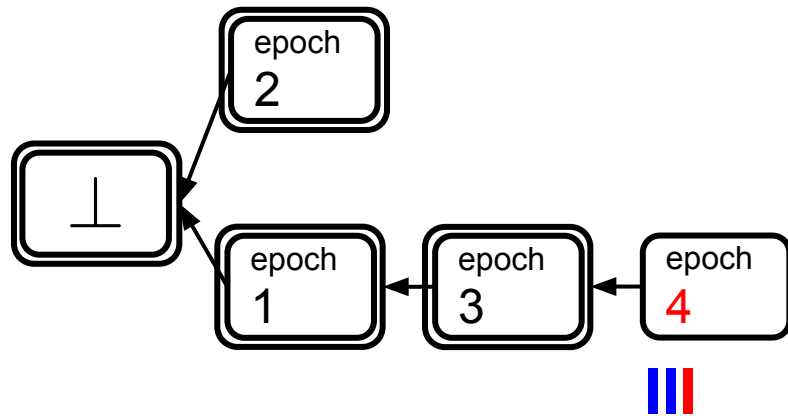
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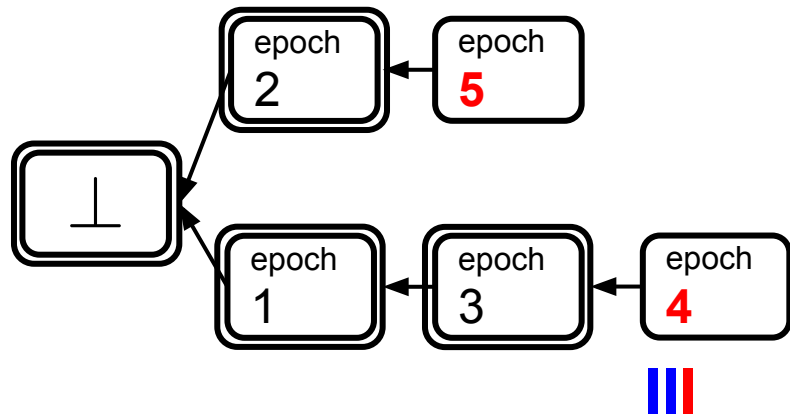
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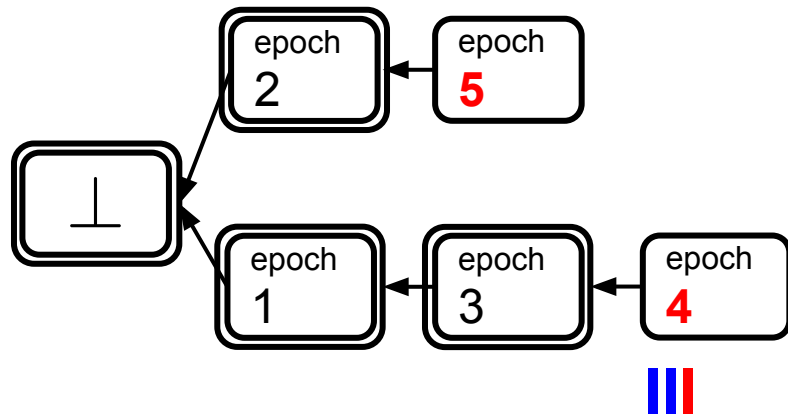
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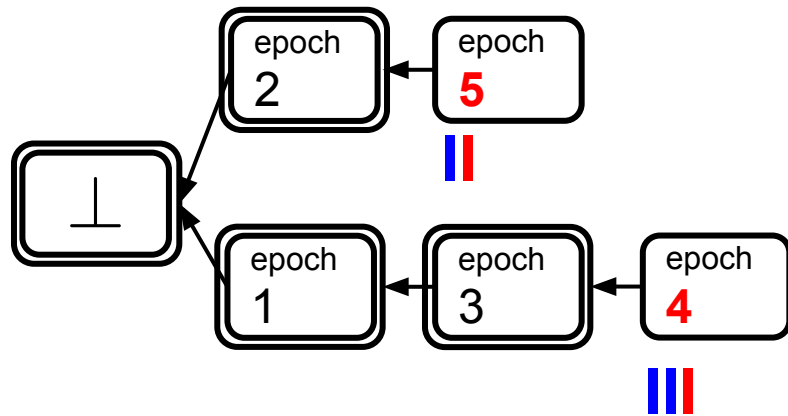
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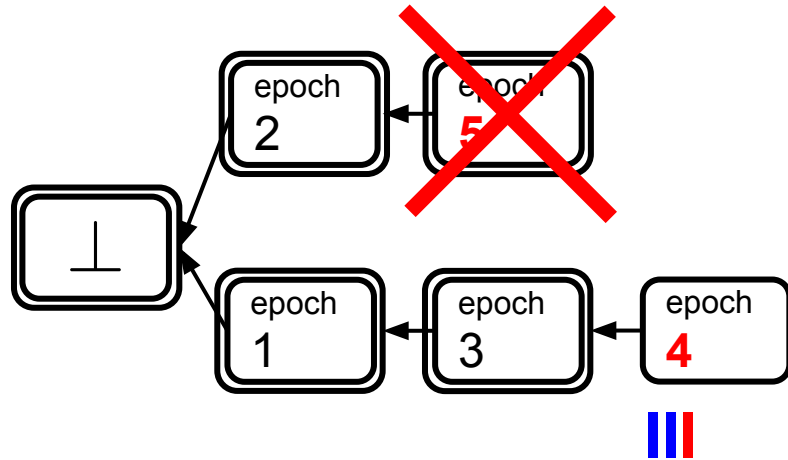
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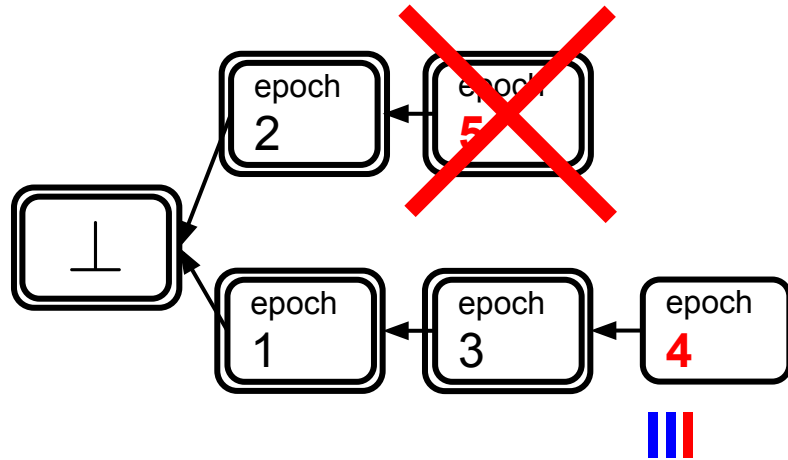
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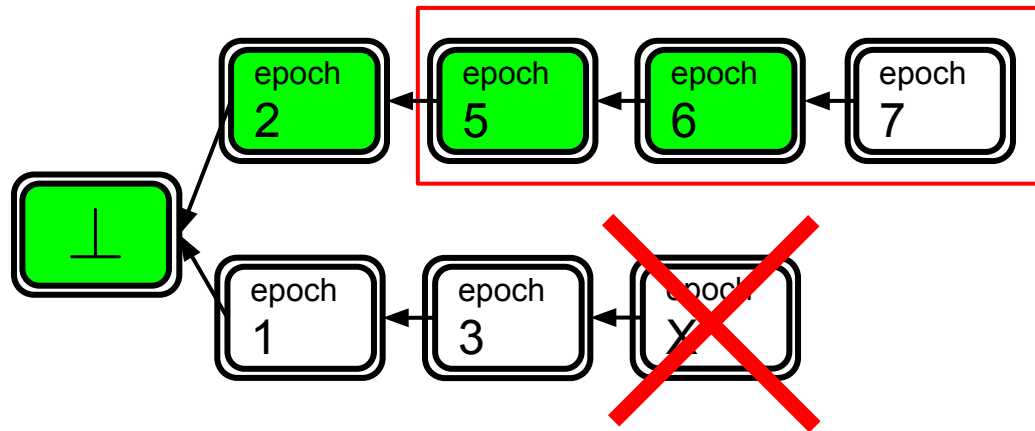
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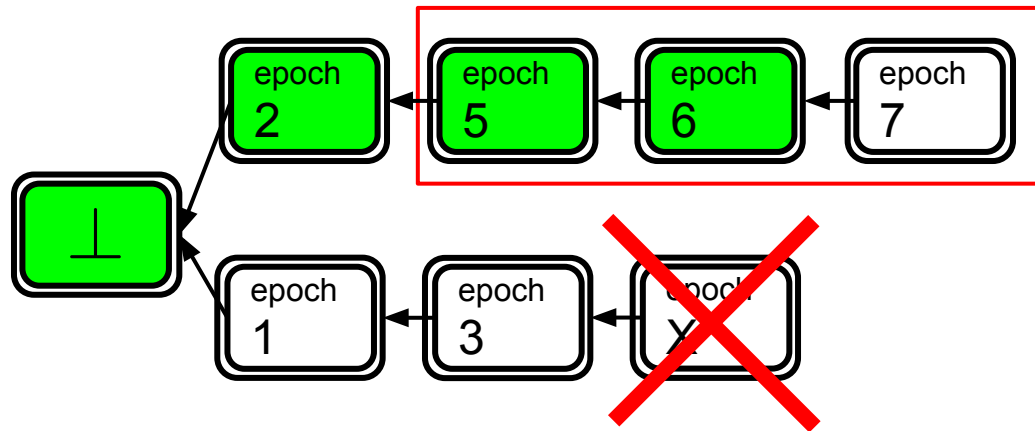
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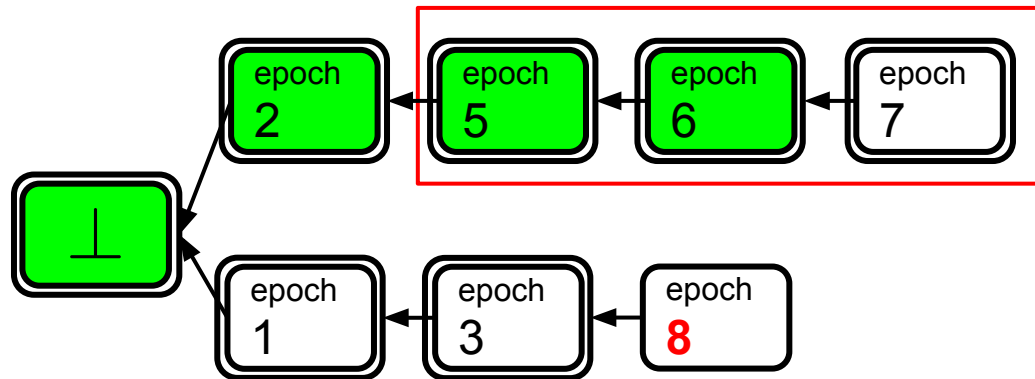
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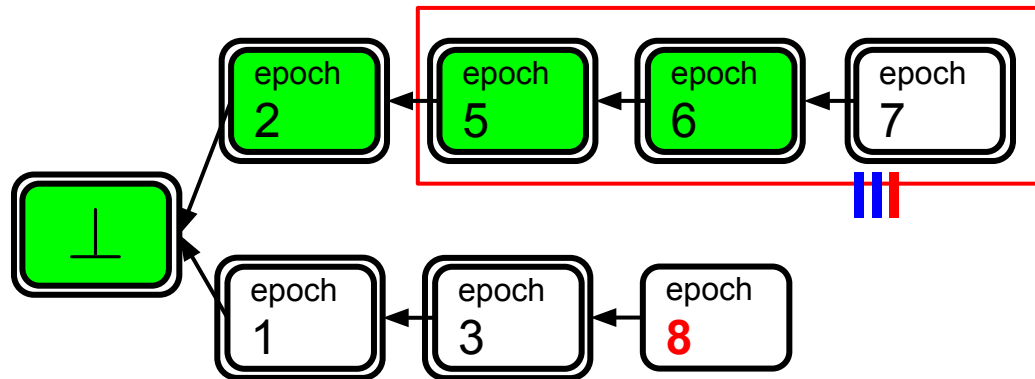
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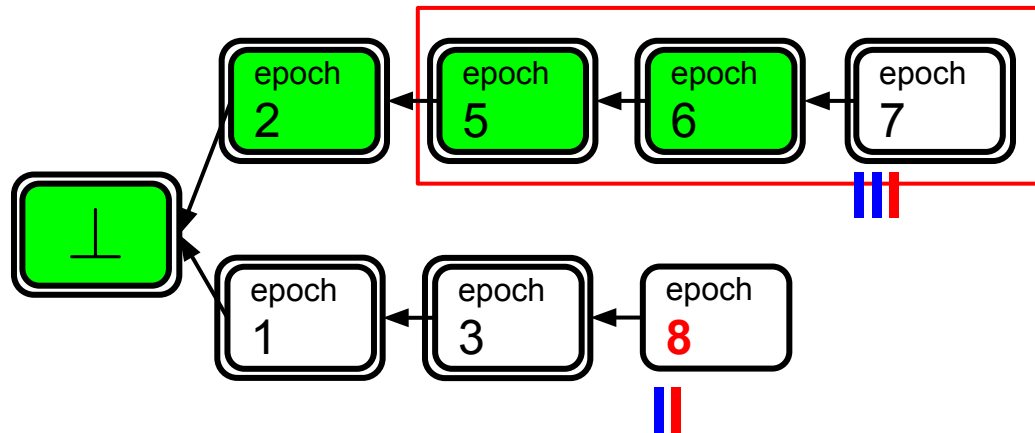
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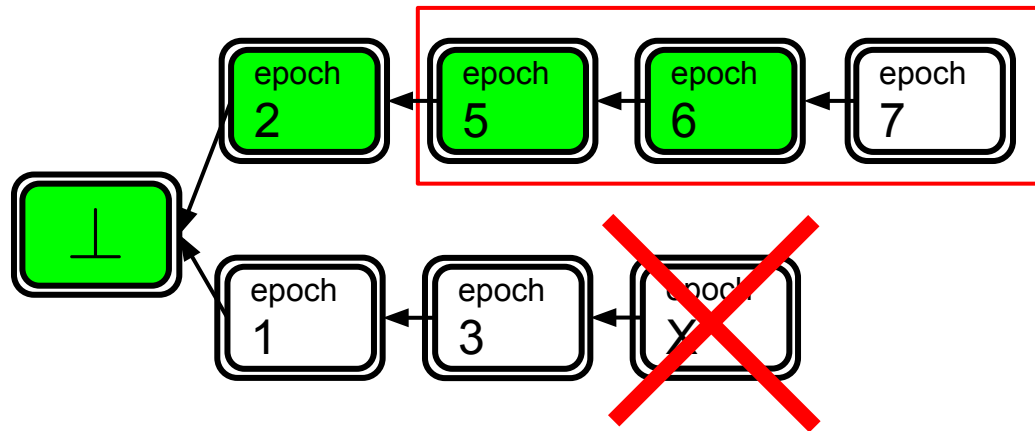
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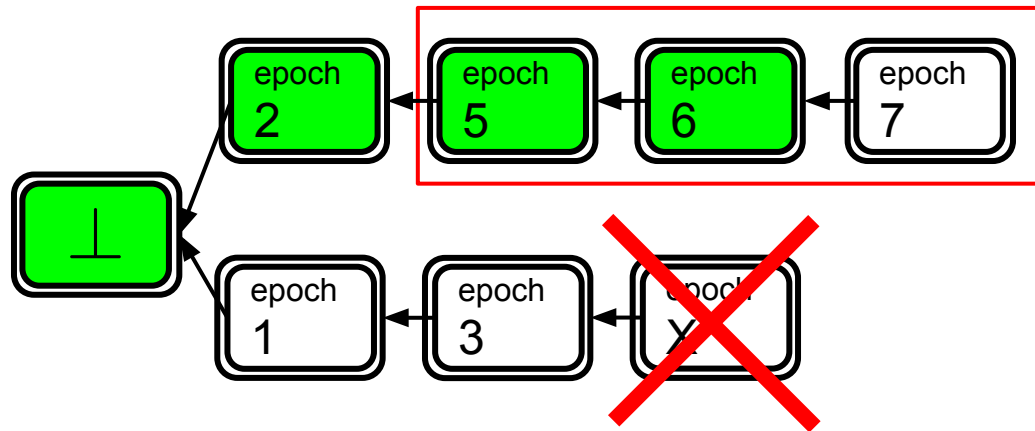
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Consistency Sketch



Intuition: Can't rewrite history

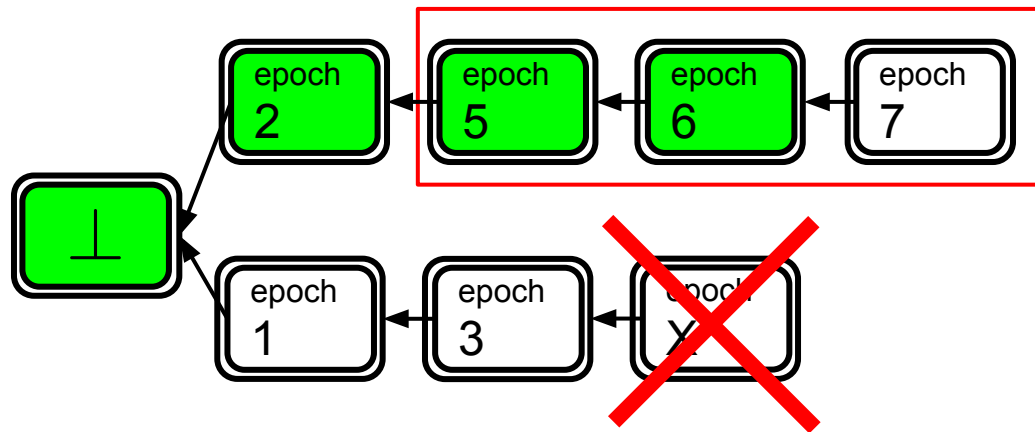
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Consistency Sketch



Intuition: Can't rewrite history

(can add a new notarized block at the same height as an existing notarized block, but never in the past)

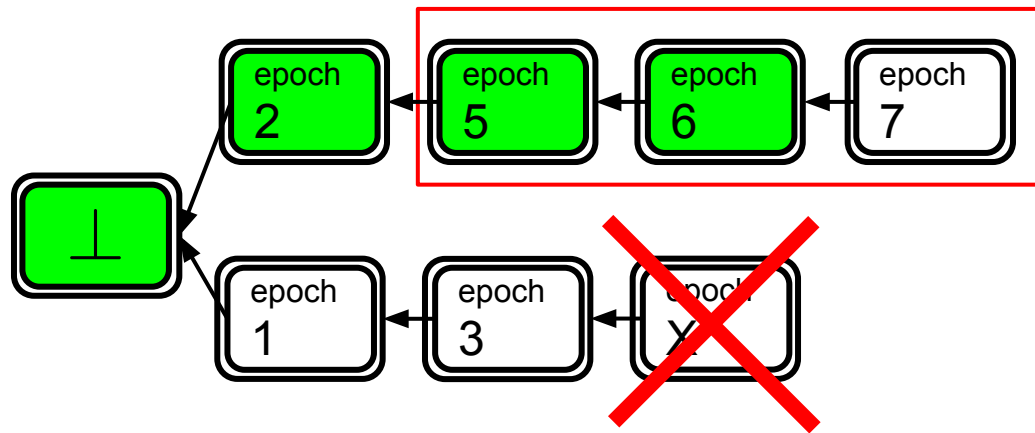
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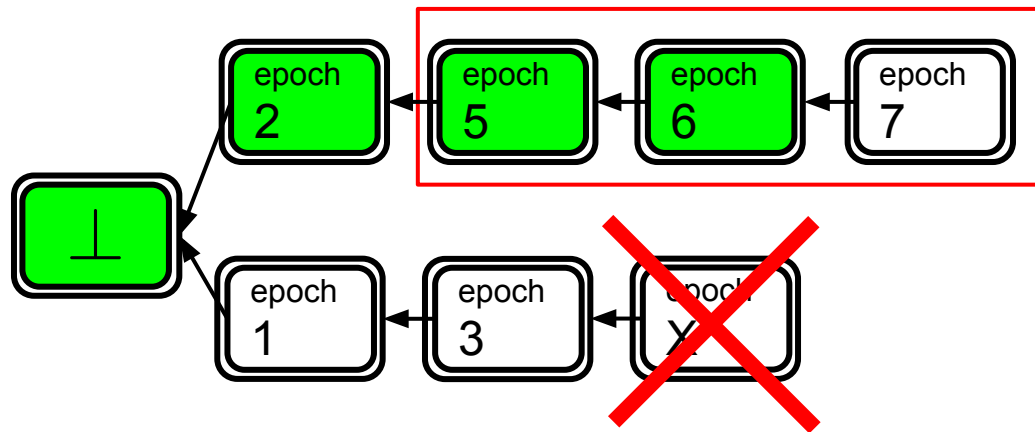
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Lemma 1: *Each epoch is associated with at most one notarized block.*

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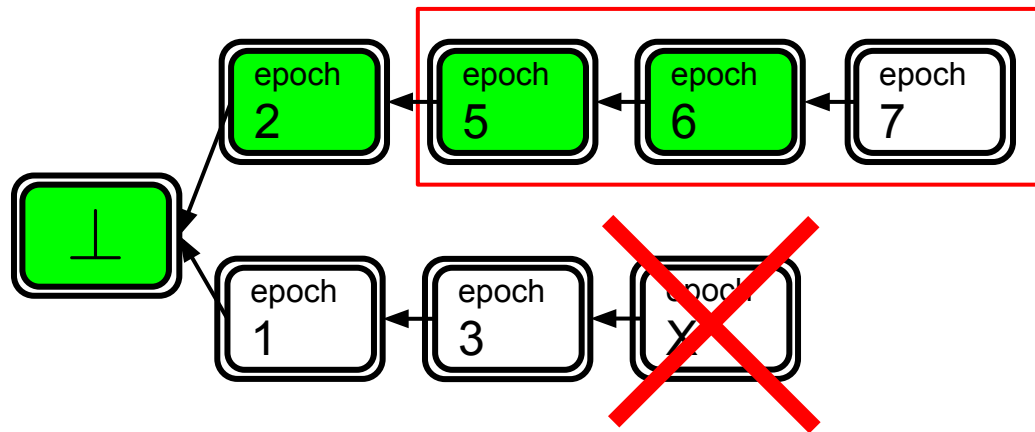


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Lemma 3: *All longer notarized chains, in any view, must extend block 6*

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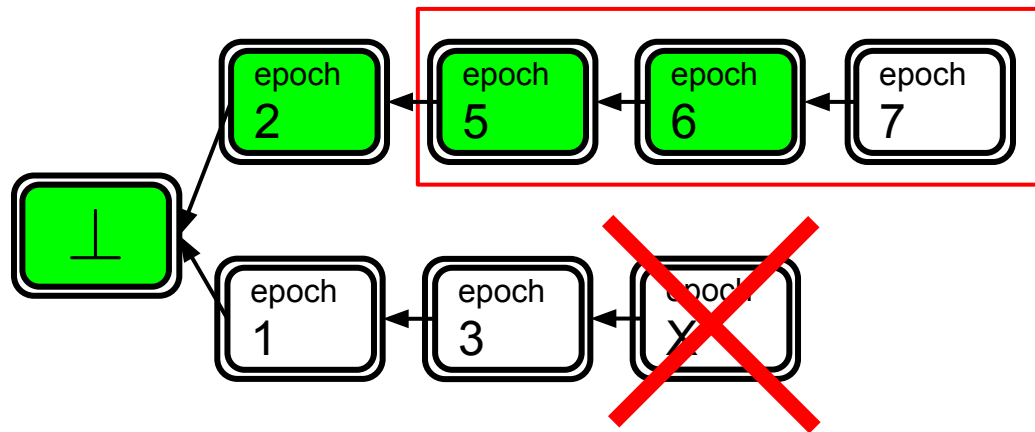


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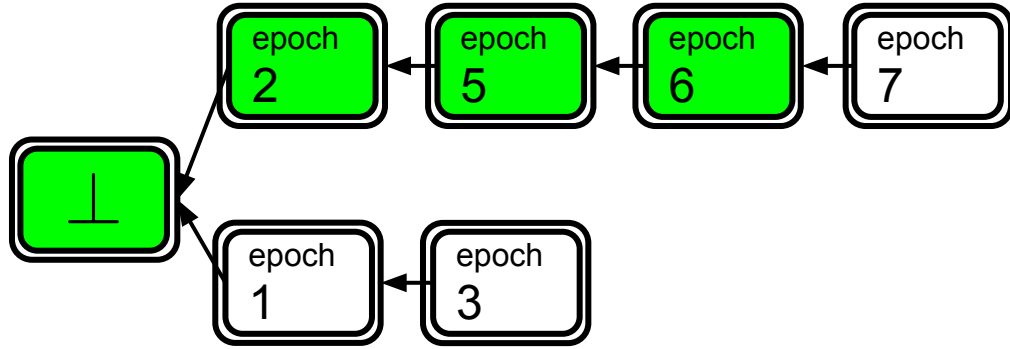
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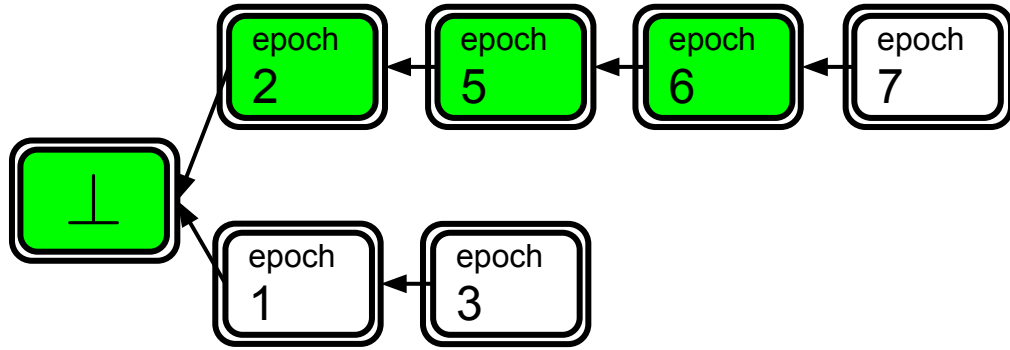
No synchrony assumptions!

Consistency Recap (Intuitive)



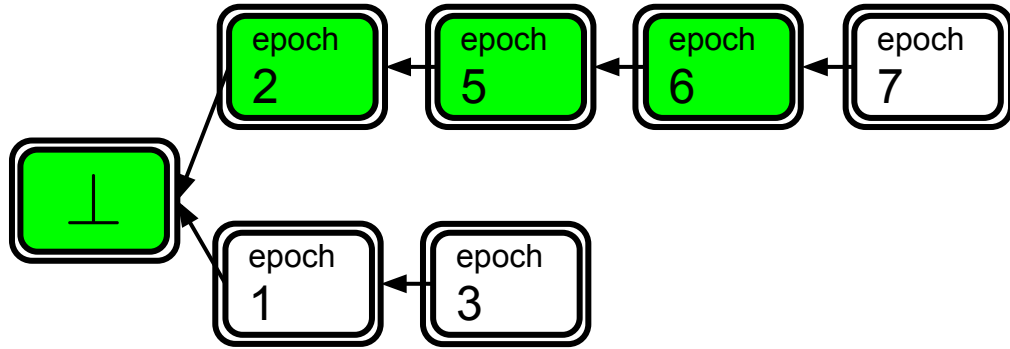
1. Can't rewrite history
2. One block per epoch
3. Demonstrate sudden chain growth

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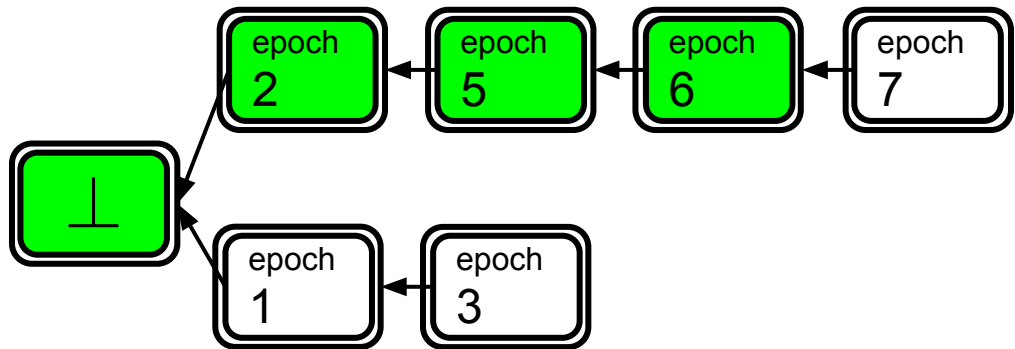


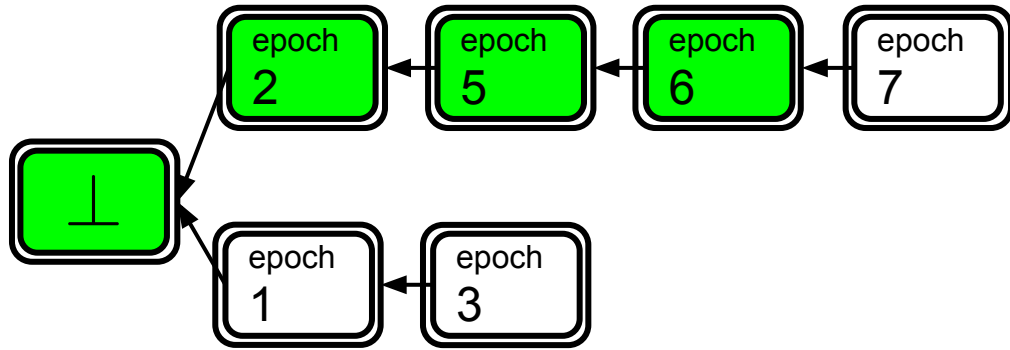
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= chain provably longer than competitors

Consistency Recap (Intuitive)



1. Can't rewrite history
2. One block per epoch
3. Demonstrate sudden chain growth (why 3?)
= chain provably longer than competitors



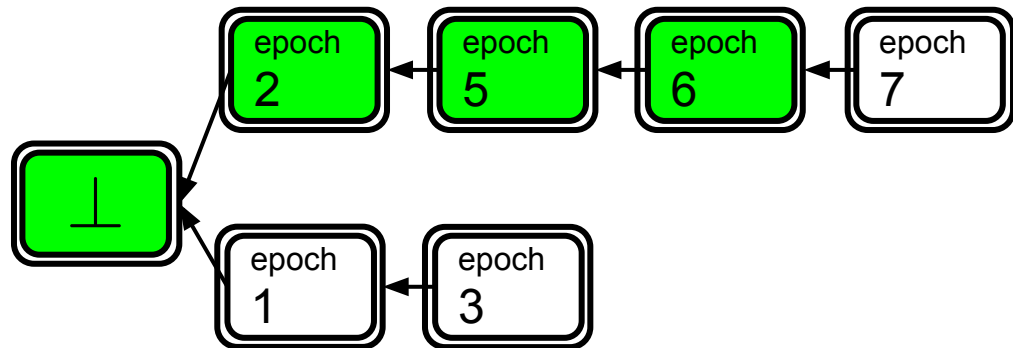


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Liveness?

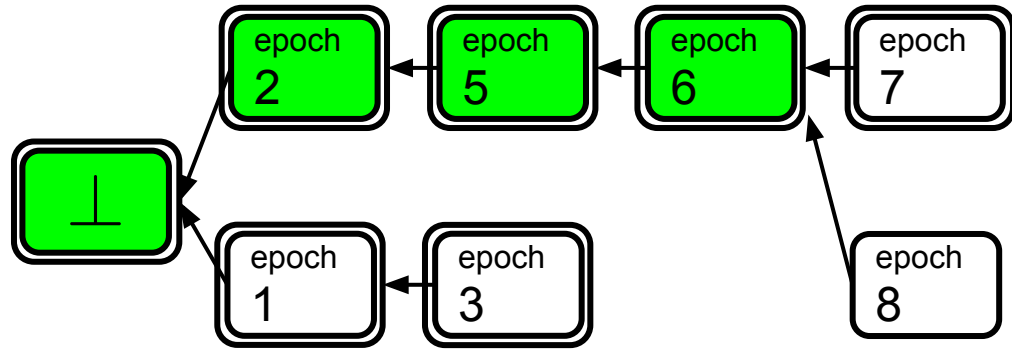


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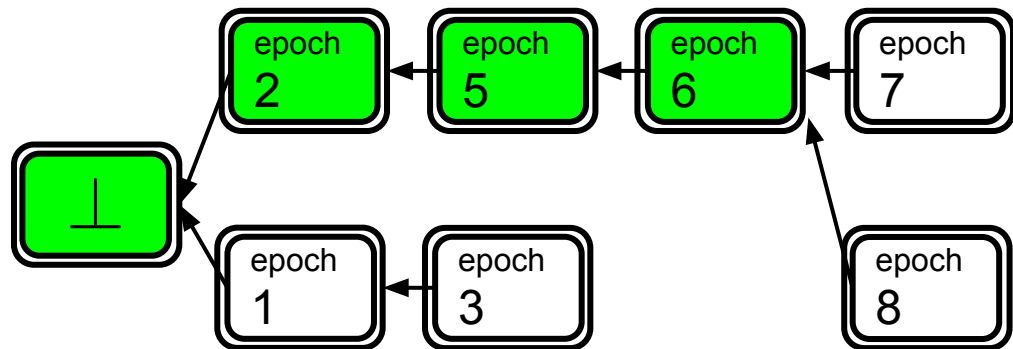


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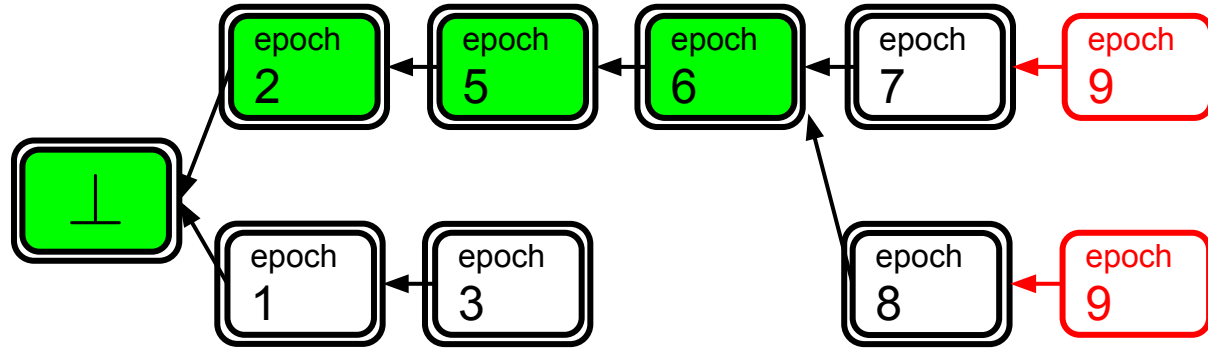


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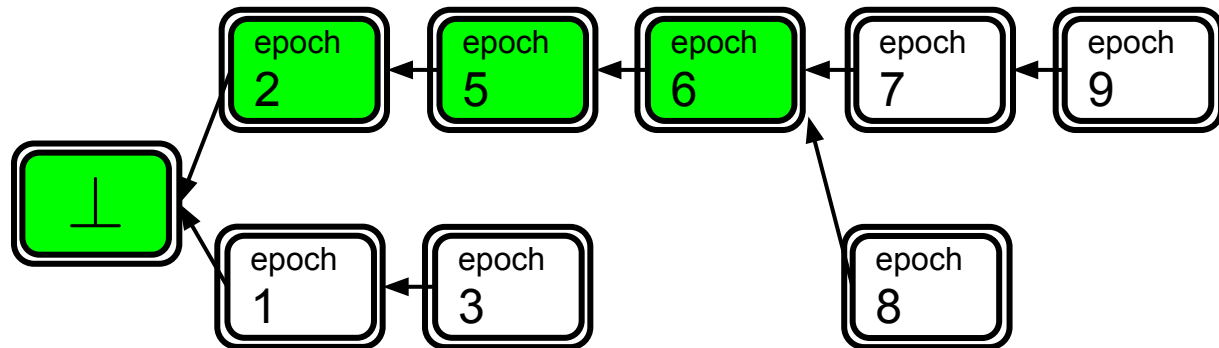


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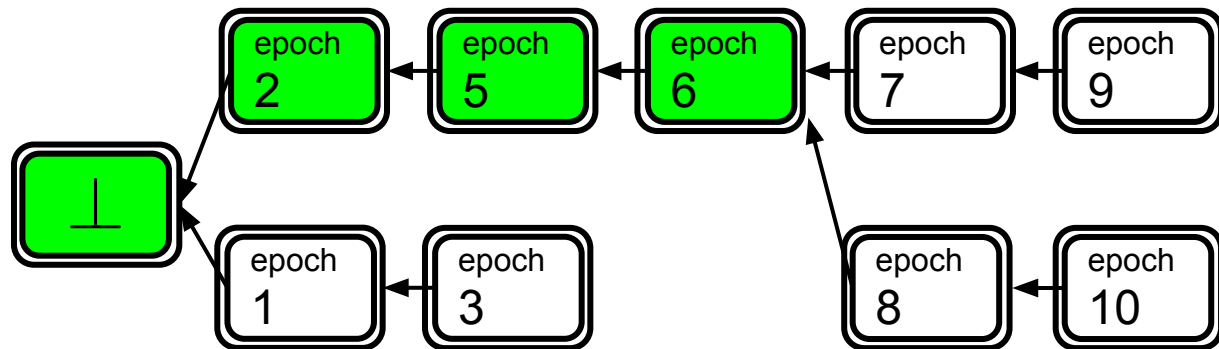


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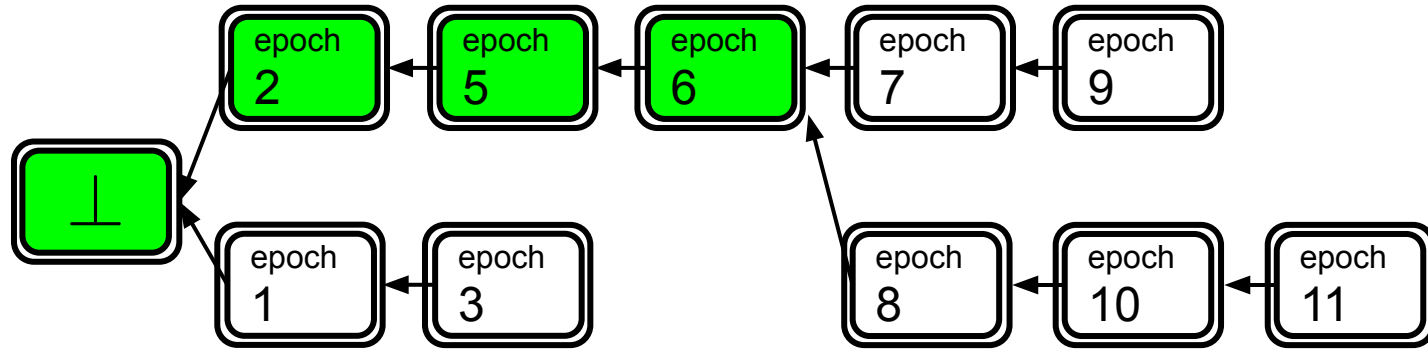


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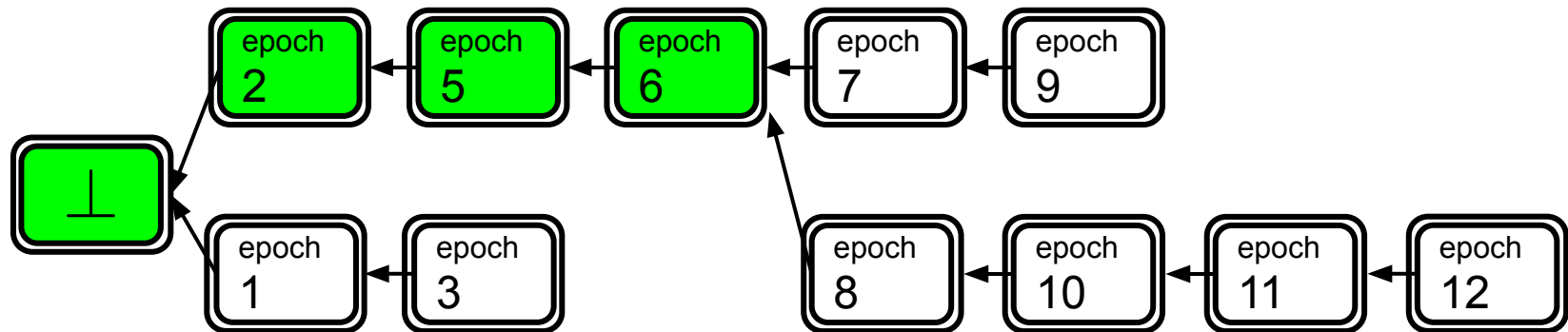


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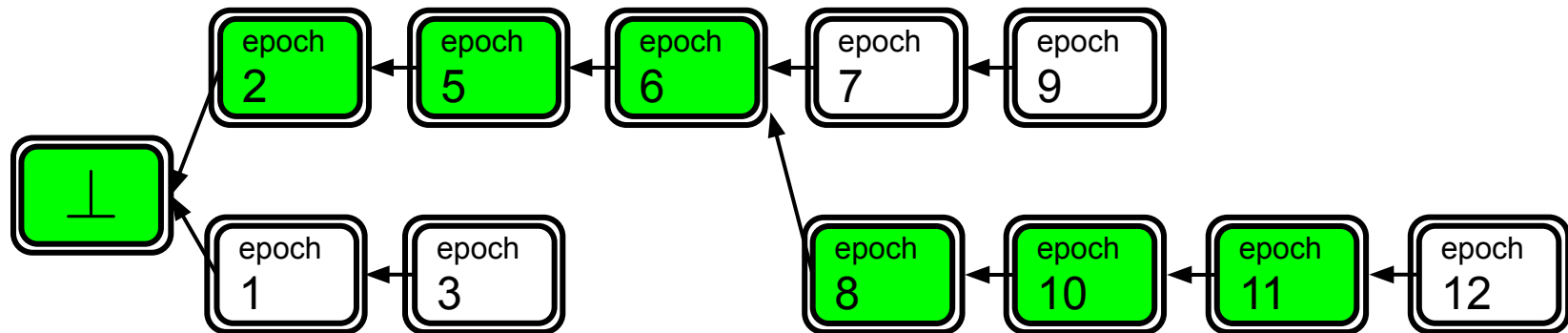


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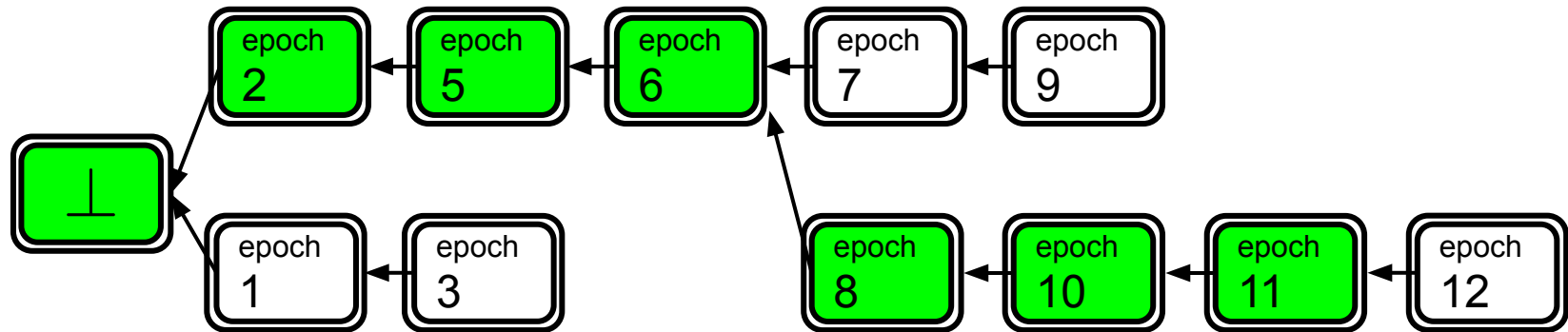


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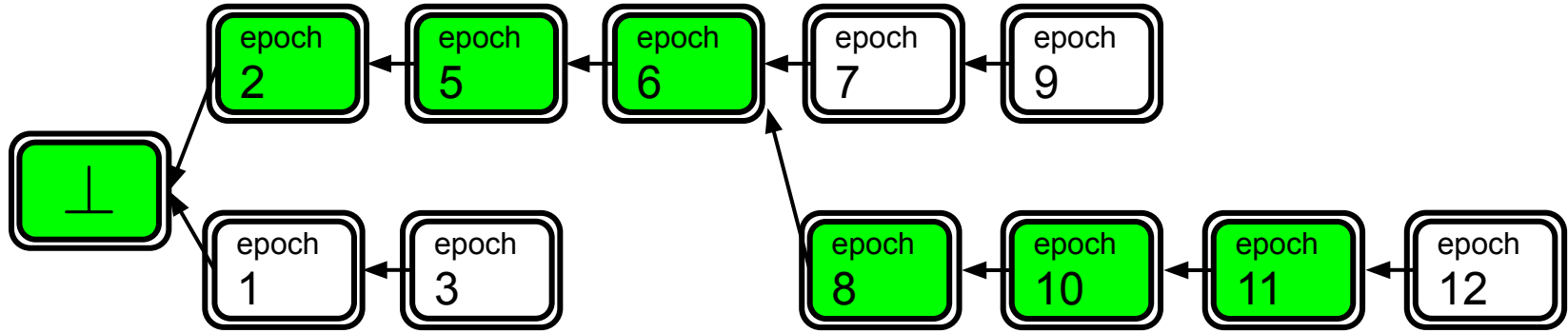
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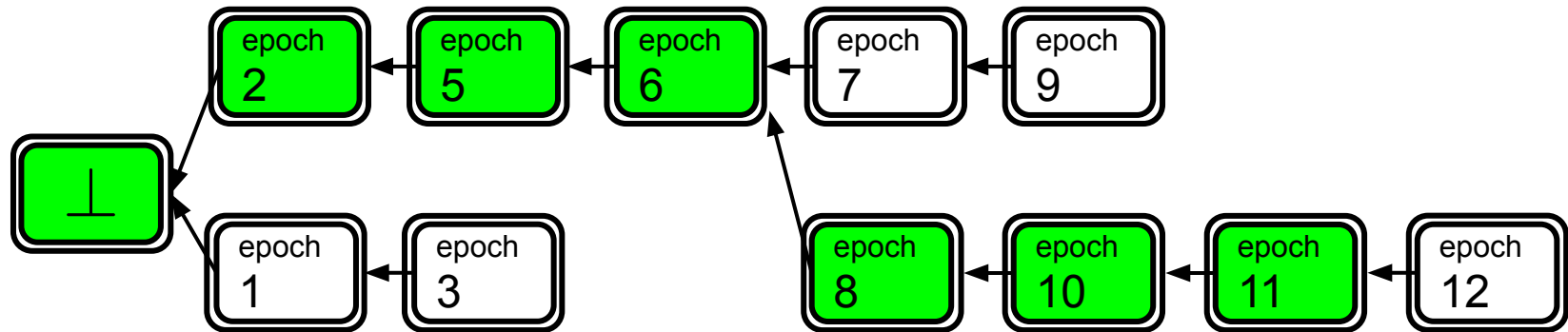
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Liveness?



We need many good leaders in a row

- Random leaders: get lucky
- Stable leader mechanism
- Not bad!



Recap

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Eprint: ia.cr/2020/088

Questions?

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Eprint: ia.cr/2020/088